Instruction Set Nomenclature:

Status Register (SREG):
SREG: Status register
C: Carry flag in status register
Z: Zero flag in status register
N: Negative flag in status register
V: Twos complement overflow indicator
S: N ⊕ V, For signed tests
H: Half Carry flag in the status register
T: Transfer bit used by BLD and BST instructions
I: Global interrupt enable/disable flag

Registers and operands:
Rd: Destination (and source) register in the register file
Rr: Source register in the register file
R: Result after instruction is executed
K: Constant literal or byte data (8 bit)
k: Constant address data for program counter
b: Bit in the register file or I/O register (3 bit)
s: Bit in the status register (3 bit)

X,Y,Z: Indirect address register (X=R27:R26, Y=R29:R28 and Z=R31:R30)
P: I/O port address
q: Displacement for direct addressing (6 bit)

I/O Registers
RAMPX, RAMPY, RAMPZ: Registers concatenated with the X, Y and Z registers enabling indirect addressing of the whole SRAM area on MCUs with more than 64K bytes SRAM.

Stack:
STACK: Stack for return address and pushed registers
SP: Stack Pointer to STACK

Flags:
⇔: Flag affected by instruction
0: Flag cleared by instruction
1: Flag set by instruction
-: Flag not affected by instruction

Conditional Branch Summary

Test	Boolean	Mnemonic	Complementary	Boolean	Mnemonic	Comment
Rd > Rr	Z•(N ⊕ V) = 0	BRLT*	Rd ≤ Rr	Z+(N ⊕ V) = 1	BRGE*	Signed
Rd ≥ Rr	(N ⊕ V) = 0	BRGE	Rd < Rr	(N ⊕ V) = 1	BRLT	Signed
Rd = Rr	Z = 1	BREQ	Rd ≠ Rr	Z = 0	BRNE	Signed
Rd ≤ Rr	Z+(N ⊕ V) = 1	BRGE*	Rd > Rr	Z•(N ⊕ V) = 0	BRLT*	Signed
Rd < Rr	(N ⊕ V) = 1	BRLT	Rd ≥ Rr	(N ⊕ V) = 0	BRGE	Signed
Rd > Rr	C + Z = 0	BRLO*	Rd ≤ Rr	C + Z = 1	BRSH*	Unsigned
Rd ≥ Rr	C = 0	BRSH/BRCC	Rd < Rr	C = 1	BRLO/BRCS	Unsigned
Rd = Rr	Z = 1	BREQ	Rd ≠ Rr	Z = 0	BRNE	Unsigned
Rd ≤ Rr	C + Z = 1	BRSH*	Rd > Rr	C + Z = 0	BRLO*	Unsigned
Rd < Rr	C = 1	BRLO/BRCS	Rd ≥ Rr	C = 0	BRSH/BRCC	Unsigned
Carry	C = 1	BRCS	No carry	C = 0	BRCC	Simple
Negative	N = 1	BRMI	Positive	N = 0	BRPL	Simple
Overflow	V = 1	BRVS	No overflow	V = 0	BRVC	Simple
Zero	Z = 1	BREQ	Not zero	Z = 0	BRNE	Simple

* Interchange Rd and Rr in the operation before the test. i.e. CP Rd,Rr → CP Rr,Rd

Complete Instruction Set Summary

Mnemonics	Operands	Description	Operation	Flags	#Clock Note
ARITHMETIC AND LOGIC INSTRUCTIONS					
ADD	Rd, Rr	Add without Carry	Rd ← Rd + Rr	Z,C,N,V,H	1
ADC	Rd, Rr	Add with Carry	Rd ← Rd + Rr + C	Z,C,N,V,H	1
ADIW	Rd, K	Add Immediate to Word	Rd+1:Rd ← Rd+1:Rd + K	Z,C,N,V	2
SUB	Rd, Rr	Subtract without Carry	Rd ← Rd - Rr	Z,C,N,V,H	1
SUBI	Rd, K	Subtract Immediate	Rd ← Rd - K	Z,C,N,V,H	1
SBC	Rd, Rr	Subtract with Carry	Rd ← Rd - Rr - C	Z,C,N,V,H	1
SBCI	Rd, K	Subtract Immediate with Carry	Rd ← Rd - K - C	Z,C,N,V,H	1
SBIW	Rd, K	Subtract Immediate from Word	Rd+1:Rd ← Rd+1:Rd - K	Z,C,N,V	2
AND	Rd, Rr	Logical AND	Rd ← Rd . Rr	Z,N,V	1
ANDI	Rd, K	Logical AND with Immediate	Rd ← Rd • K	Z,N,V	1
OR	Rd, Rr	Logical OR	Rd ← Rd v Rr	Z,N,V	1
ORI	Rd, K	Logical OR with Immediate	Rd ← Rd v K	Z,N,V	1
EOR	Rd, Rr	Exclusive OR	Rd ← Rd ⊕ Rr	Z,N,V	1
COM	Rd	One's Complement	Rd ← $FF - Rd	Z,C,N,V	1
NEG	Rd	Two's Complement	Rd ← $00 - Rd	Z,C,N,V,H	1
SBR	Rd,K	Set Bit(s) in Register	Rd ← Rd v K	Z,N,V	1
CBR	Rd,K	Clear Bit(s) in Register	Rd ← Rd • ($FFh - K)	Z,N,V	1
INC	Rd	Increment	Rd ← Rd + 1	Z,N,V	1
DEC	Rd	Decrement	Rd ← Rd - 1	Z,N,V	1
TST	Rd	Test for Zero or Minus	Rd ← Rd • Rd	Z,N,V	1
CLR	Rd	Clear Register	Rd ← Rd ⊕ Rd	Z,N,V	1
SER	Rd	Set Register	Rd ← $FF	None	1
CP	Rd,Rr	Compare	Rd - Rr	Z,C,N,V,H,	1
CPC	Rd,Rr	Compare with Carry	Rd - Rr - C	Z,C,N,V,H	1
CPI	Rd,K	Compare with Immediate	Rd - K	Z,C,N,V,H	1

√) Not available in base-line microcontrollers

(continued)

Complete Instruction Set Summary (continued)

Mnemonics	Operands	Description	Operation	Flags	#Clock Note
BRANCH INSTRUCTIONS					
RJMP	k	Relative Jump	PC ← PC + k + 1	None	2
IJMP		Indirect Jump to (Z)	PC ← Z	None	2
JMP	k	Jump	PC ← k	None	3
RCALL	k	Relative Call Subroutine	PC ← PC + k + 1	None	3
ICALL		Indirect Call to (Z)	PC ← Z	None	3
CALL	k	Call Subroutine	PC ← k	None	4
RET		Subroutine Return	PC ← STACK	None	4
RETI		Interrupt Return	PC ← STACK	I	4
CPSE	Rd,Rr	Compare, Skip if Equal	if (Rd = Rr) PC ← PC + 2 or 3	None	1 / 2 / 3
SBRC	Rr, b	Skip if Bit in Register Cleared	if (Rr(b)=0) PC ← PC + 2 or 3	None	1 / 2 / 3
SBRS	Rr, b	Skip if Bit in Register Set	if (Rr(b)=1) PC ← PC + 2 or 3	None	1 / 2 / 3
SBIC	P, b	Skip if Bit in I/O Register Cleared	if(I/O(P,b)=0) PC ← PC + 2 or 3	None	1 / 2 / 3
SBIS	P, b	Skip if Bit in I/O Register Set	If(I/O(P,b)=1) PC← PC + 2 or 3	None	1 / 2 / 3
BRBS	s, k	Branch if Status Flag Set	if (SREG(s) = 1) then PC←PC+k + 1	None	1 / 2
BRBC	s, k	Branch if Status Flag Cleared	if (SREG(s) = 0) then PC←PC+k + 1	None	1 / 2
BREQ	k	Branch if Equal	if (Z = 1) then PC ← PC + k + 1	None	1 / 2
BRNE	k	Branch if Not Equal	if (Z = 0) then PC ← PC + k + 1	None	1 / 2
BRCS	k	Branch if Carry Set	if (C = 1) then PC ← PC + k + 1	None	1 / 2
BRCC	k	Branch if Carry Cleared	if (C = 0) then PC ← PC + k + 1	None	1 / 2
BRSH	k	Branch if Same or Higher	if (C = 0) then PC ← PC + k + 1	None	1 / 2
BRLO	k	Branch if Lower	if (C = 1) then PC ← PC + k + 1	None	1 / 2
BRMI	k	Branch if Minus	if (N = 1) then PC ← PC + k + 1	None	1 / 2
BRPL	k	Branch if Plus	if (N = 0) then PC ← PC + k + 1	None	1 / 2
BRGE	k	Branch if Greater or Equal, Signed	if (N ⊕ V= 0) then PC ← PC+ k + 1	None	1 / 2
BRLT	k	Branch if Less Than, Signed	if (N ⊕ V= 1) then PC ← PC + k + 1	None	1 / 2
BRHS	k	Branch if Half Carry Flag Set	if (H = 1) then PC ← PC + k + 1	None	1 / 2
BRHC	k	Branch if Half Carry Flag Cleared	if (H = 0) then PC ← PC + k + 1	None	1 / 2
BRTS	k	Branch if T Flag Set	if (T = 1) then PC ← PC + k + 1	None	1 / 2
BRTC	k	Branch if T Flag Cleared	if (T = 0) then PC ← PC + k + 1	None	1 / 2
BRVS	k	Branch if Overflow Flag is Set	if (V = 1) then PC ← PC + k + 1	None	1 / 2
BRVC	k	Branch if Overflow Flag is Cleared	if (V = 0) then PC ← PC + k + 1	None	1 / 2
BRIE	k	Branch if Interrupt Enabled	if (I = 1) then PC ← PC + k + 1	None	1 / 2
BRID	k	Branch if Interrupt Disabled	if (I = 0) then PC ← PC + k + 1	None	1 / 2

(continued)

Complete Instruction Set Summary (continued)

Mnemonics	Operands	Description	Operation	Flags	#Clock Note
DATA TRANSFER INSTRUCTIONS					
MOV	Rd, Rr	Copy Register	Rd ← Rr	None	1
LDI	Rd, K	Load Immediate	Rd ← K	None	1
LDS	Rd, k	Load Direct from SRAM	Rd ← (k)	None	3
LD	Rd, X	Load Indirect	Rd ← (X)	None	2
LD	Rd, X+	Load Indirect and Post-Increment	Rd ← (X), X ← X + 1	None	2
LD	Rd, -X	Load Indirect and Pre-Decrement	X ← X - 1, Rd ← (X)	None	2
LD	Rd, Y	Load Indirect	Rd ← (Y)	None	2
LD	Rd, Y+	Load Indirect and Post-Increment	Rd ← (Y), Y ← Y + 1	None	2
LD	Rd, -Y	Load Indirect and Pre-Decrement	Y ← Y - 1, Rd ← (Y)	None	2
LDD	Rd,Y+q	Load Indirect with Displacement	Rd ← (Y + q)	None	2
LD	Rd, Z	Load Indirect	Rd ← (Z)	None	2
LD	Rd, Z+	Load Indirect and Post-Increment	Rd ← (Z), Z ← Z+1	None	2
LD	Rd, -Z	Load Indirect and Pre-Decrement	Z ← Z - 1, Rd ← (Z)	None	2
LDD	Rd, Z+q	Load Indirect with Displacement	Rd ← (Z + q)	None	2
STS	k, Rr	Store Direct to SRAM	Rd ← (k)	None	3
ST	X, Rr	Store Indirect	(X) ← Rr	None	2
ST	X+, Rr	Store Indirect and Post-Increment	(X) ← Rr, X ← X + 1	None	2
ST	-X, Rr	Store Indirect and Pre-Decrement	X ← X - 1, (X) ← Rr	None	2
ST	Y, Rr	Store Indirect	(Y) ← Rr	None	2
ST	Y+, Rr	Store Indirect and Post-Increment	(Y) ← Rr, Y ← Y + 1	None	2
ST	-Y, Rr	Store Indirect and Pre-Decrement	Y ← Y - 1, (Y) ← Rr	None	2
STD	Y+q,Rr	Store Indirect with Displacement	(Y + q) ← Rr	None	2
ST	Z, Rr	Store Indirect	(Z) ← Rr	None	2
ST	Z+, Rr	Store Indirect and Post-Increment	(Z) ← Rr, Z ← Z + 1	None	2
ST	-Z, Rr	Store Indirect and Pre-Decrement	Z ← Z - 1, (Z) ← Rr	None	2
STD	Z+q,Rr	Store Indirect with Displacement	(Z + q) ← Rr	None	2
LPM		Load Program Memory	R0 ← (Z)	None	3
IN	Rd, P	In Port	Rd ← P	None	1
OUT	P, Rr	Out Port	P ← Rr	None	1
PUSH	Rr	Push Register on Stack	STACK ← Rr	None	2
POP	Rd	Pop Register from Stack	Rd ← STACK	None	2

(continued)

Complete Instruction Set Summary (continued)

Mnemonics	Operands	Description	Operation	Flags	#Clock Note
BIT AND BIT-TEST INSTRUCTIONS					
LSL	Rd	Logical Shift Left	Rd(n+1)←Rd(n),Rd(0)←0,C←Rd(7	Z,C,N,V,H	1
LSR	Rd	Logical Shift Right	Rd(n)←Rd(n+1),Rd(7)←0,C←Rd(0	Z,C,N,V	1
ROL	Rd	Rotate Left Through Carry	Rd(0)←C,Rd(n+1)←Rd(n),C←Rd(7	Z,C,N,V,H	1
ROR	Rd	Rotate Right Through Carry	Rd(7)←C,Rd(n)←Rd(n+1),C←Rd(0	Z,C,N,V	1
ASR	Rd	Arithmetic Shift Right	Rd(n) ← Rd(n+1), n=0..6	Z,C,N,V	1
SWAP	Rd	Swap Nibbles	Rd(3..0) ↔ Rd(7..4)	None	1
BSET	s	Flag Set	SREG(s) ← 1	SREG(s)	1
BCLR	s	Flag Clear	SREG(s) ← 0	SREG(s)	1
SBI	P, b	Set Bit in I/O Register	I/O(P, b) ← 1	None	2
CBI	P, b	Clear Bit in I/O Register	I/O(P, b) ← 0	None	2
BST	Rr, b	Bit Store from Register to T	T ← Rr(b)	T	1
BLD	Rd, b	Bit load from T to Register	Rd(b) ← T	None	1
SEC		Set Carry	C ← 1	C	1
CLC		Clear Carry	C ← 0	C	1
SEN		Set Negative Flag	N ← 1	N	1
CLN		Clear Negative Flag	N ← 0	N	1
SEZ		Set Zero Flag	Z ← 1	Z	1
CLZ		Clear Zero Flag	Z ← 0	Z	1
SEI		Global Interrupt Enable	I ← 1	I	1
CLI		Global Interrupt Disable	I ← 0	I	1
SES		Set Signed Test Flag	S ← 1	S	1
CLS		Clear Signed Test Flag	S ← 0	S	1
SEV		Set Two’s Complement Overflow	V ← 1	V	1
CLV		Clear Two’s Complement Overflow	V ← 0	V	1
SET		Set T in SREG	T ← 1	T	1
CLT		Clear T in SREG	T ← 0	T	1
SEH		Set Half Carry Flag in SREG	H ← 1	H	1
CLH		Clear Half Carry Flag in SREG	H ← 0	H	1
NOP		No Operation		None	1
SLEEP		Sleep	(see specific descr. for Sleep)	None	1
WDR		Watchdog Reset	(see specific descr. for WDR)	None	1

ADC - Add with Carry

Description:
Adds two registers and the contents of the C flag and places the result in the destination register Rd.

Operation:
(i) Rd ← Rd + Rr + C

	Syntax:	Operands:	Program Counter:
(i)	ADC Rd,Rr	0 ≤ d ≤ 31, 0 ≤ r ≤ 31	PC ← PC + 1

16 bit Opcode:

0001	11rd	dddd	rrrr

Status Register (SREG) Boolean Formulae:

I	T	H	S	V	N	Z	C
–	–	⇔	⇔	⇔	⇔	⇔	⇔

H: Rd3•Rr3+Rr3+$\overline{R3}$+$\overline{R3}$•Rd3
Set if there was a carry from bit 3; cleared otherwise

S: N ⊕ V, For signed tests.

V: Rd7•Rr7•$\overline{R7}$+$\overline{Rd7}$•$\overline{Rr7}$•R7
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{Rd7}$• $\overline{Rr7}$• $\overline{Rr7}$ • $\overline{R7}$ • $\overline{R7}$ •$\overline{Rd7}$
Set if the result is $00; cleared otherwise.

C: Rd7•Rr7+Rr7•$\overline{R7}$+$\overline{R7}$•Rd7
Set if there was carry from the MSB of the result; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
                ; Add R1:R0 to R3:R2
add   r2,r0     ; Add low byte
adc   r3,r1     ; Add with carry high byte
```

Words: 1 (2 bytes)
Cycles: 1

ADD - Add without Carry

Description:

Adds two registers without the C flag and places the result in the destination register Rd.

Operation:

(i) Rd ← Rd + Rr

	Syntax:	Operands:	Program Counter:
(i)	ADD Rd,Rr	0 ≤ d ≤ 31, 0 ≤ r ≤ 31	PC ← PC + 1

16 bit Opcode:

0000	11rd	dddd	rrrr

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: Rd3•Rr3+Rr3+$\overline{R3}$+$\overline{R3}$•Rd3
Set if there was a carry from bit 3; cleared otherwise

S: N ⊕ V, For signed tests.

V: Rd7•Rr7•$\overline{R7}$+$\overline{Rd7}$•$\overline{Rr7}$•R7
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}$• $\overline{R6}$ •$\overline{R5}$• $\overline{R4}$ •$\overline{R3}$ •$\overline{R2}$ •$\overline{R1}$ •$\overline{R0}$
Set if the result is $00; cleared otherwise.

C: Rd7 •Rr7 +Rr7 •$\overline{R7}$+ $\overline{R7}$ •Rd7
Set if there was carry from the MSB of the result; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
add   r1,r2     ; Add r2 to r1 (r1=r1+r2)
add   r28,r28   ; Add r28 to itself (r28=r28+r28)
```

Words: 1 (2 bytes)
Cycles: 1

ADIW - Add Immediate to Word

Description:

Adds an immediate value (0-63) to a register pair and places the result in the register pair. This instruction operates on the upper four register pairs, and is well suited for operations on the pointer registers.

Operation:

(i) Rdh:Rdl ← Rdh:Rdl + K

	Syntax:	Operands:	Program Counter:
(i)	ADIW Rdl,K	dl ∈ {24,26,28,30}, 0 ≤ K ≤ 63	PC ← PC + 1

16 bit Opcode:

1001	0110	KKdd	KKKK

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	⇔	⇔	⇔	⇔

S: N ⊕ V, For signed tests.

V: Rdh7 R15
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R15
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R15} \bullet \overline{R14} \bullet \overline{R13} \bullet \overline{R12} \bullet \overline{R11} \bullet \overline{R10} \bullet \overline{R9} \bullet \overline{R8} \bullet \overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $0000; cleared otherwise.

C: $\overline{R15} \bullet$ Rdh7
Set if there was carry from the MSB of the result; cleared otherwise.

R (Result) equals Rdh:Rdl after the operation (Rdh7-Rdh0 = R15-R8, Rdl7-Rdl0=R7-R0).

Example:

```
adiw  r24,1      ; Add 1 to r25:r24
adiw  r30,63     ; Add 63 to the Z pointer(r31:r30)
```

Words: 1 (2 bytes)
Cycles: 2

AND - Logical AND

Description:

Performs the logical AND between the contents of register Rd and register Rr and places the result in the destination register Rd.

Operation:

(i) Rd ← Rd • Rr

	Syntax:	Operands:	Program Counter:
(i)	AND Rd,Rr	0 ≤ d ≤ 31, 0 ≤ r ≤ 31	PC ← PC + 1

16 bit Opcode:

0010	00rd	dddd	rrrr

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	**0**	⇔	⇔	-

S: N ⊕ V, For signed tests.

V: 0
Cleared

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
and   r2,r3      ; Bitwise and r2 and r3, result in r2
ldi   r16,1      ; Set bitmask 0000 0001 in r16
and   r2,r16     ; Isolate bit 0 in r2
```

Words: 1 (2 bytes)
Cycles: 1

ANDI - Logical AND with Immediate

Description:

Performs the logical AND between the contents of register Rd and a constant and places the result in the destination register Rd.

Operation:

(i) Rd ← Rd • K

	Syntax:	Operands:	Program Counter:
(i)	ANDI Rd,K	16 ≤ d ≤ 31, 0 ≤ K ≤ 255	PC ← PC + 1

16 bit Opcode:

0111	KKKK	dddd	KKKK

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	**0**	⇔	⇔	-

S: N ⊕ V, For signed tests.

V: 0
Cleared

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}$ •$\overline{R6}$• $\overline{R5}$•$\overline{R4}$ •$\overline{R3}$• $\overline{R2}$• $\overline{R1}$• $\overline{R0}$
Set if the result is $00; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
andi  r17,$0F    ; Clear upper nibble of r17
andi  r18,$10    ; Isolate bit 4 in r18
andi  r19,$AA    ; Clear odd bits of r19
```

Words: 1 (2 bytes)
Cycles: 1

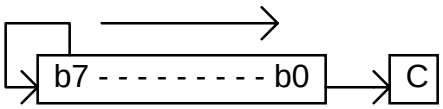
ASR - Arithmetic Shift Right

Description:

Shifts all bits in Rd one place to the right. Bit 7 is held constant. Bit 0 is loaded into the C flag of the SREG. This operation effectively divides a twos complement value by two without changing its sign. The carry flag can be used to round the result.

Operation:

(i)



	Syntax:	Operands:	Program Counter:
(i)	ASR Rd	$0 \le d \le 31$	PC ← PC + 1

16 bit Opcode:

1001	010d	dddd	0101

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	⇔	⇔	⇔	⇔

S: N ⊕ V, For signed tests.

V: N ⊕ C (For N and C after the shift)
Set if (N is set and C is clear) or (N is clear and C is set); Cleared otherwise (for values of N and C after the shift).

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; cleared otherwise.

C: Rd0
Set if, before the shift, the LSB of Rd was set; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
ldi   r16,$10      ; Load decimal 16 into r16
asr   r16          ; r16=r16 / 2
ldi   r17,$FC      ; Load -4 in r17
asr   r17          ; r17=r17/2
```

Words: 1 (2 bytes)
Cycles: 1

BCLR - Bit Clear in SREG

Description:
Clears a single flag in SREG.

Operation:
(i) SREG(s) ← 0

	Syntax:	Operands:	Program Counter:
(i)	BCLR s	0 ≤ s ≤ 7	PC ← PC + 1

16 bit Opcode:

1001	0100	1sss	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
⇔	⇔	⇔	⇔	⇔	⇔	⇔	⇔

I: 0 if s = 7; Unchanged otherwise.

T: 0 if s = 6; Unchanged otherwise.

H: 0 if s = 5; Unchanged otherwise.

S: 0 if s = 4; Unchanged otherwise.

V: 0 if s = 3; Unchanged otherwise.

N: 0 if s = 2; Unchanged otherwise.

Z: 0 if s = 1; Unchanged otherwise.

C: 0 if s = 0; Unchanged otherwise.

Example:
```
bclr    0       ; Clear carry flag
bclr    7       ; Disable interrupts
```

Words: 1 (2 bytes)
Cycles: 1

BLD - Bit Load from the T Flag in SREG to a Bit in Register.

Description:
Copies the T flag in the SREG (status register) to bit b in register Rd.

Operation:
(i) Rd(b) ← T

	Syntax:	Operands:	Program Counter:
(i)	BLD Rd,b	0 ≤ d ≤ 31, 0 ≤ b ≤ 7	PC ← PC + 1

16 bit Opcode:

1111	100d	dddd	0bbb

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
                ; Copy bit
    bst   r1,2  ; Store bit 2 of r1 in T flag
    bld   r0,4  ; Load T flag into bit 4 of r0
```

Words: 1 (2 bytes)
Cycles: 1

BRBC - Branch if Bit in SREG is Cleared

Description:

Conditional relative branch. Tests a single bit in SREG and branches relatively to PC if the bit is cleared. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form.

Operation:

(i) If SREG(s) = 0 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRBC s,k	0 ≤ s ≤ 7, -64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	01kk	kkkk	ksss

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
cpi   r20,5     ; Compare r20 to the value 5
brbc  1,noteq   ; Branch if zero flag cleared
...
noteq:nop       ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRBS - Branch if Bit in SREG is Set

Description:

Conditional relative branch. Tests a single bit in SREG and branches relatively to PC if the bit is set. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form.

Operation:

(i) If SREG(s) = 1 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRBS s,k	0 ≤ s ≤ 7, -64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	00kk	kkkk	ksss

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
          bst   r0,3     ; Load T bit with bit 3 of r0
          brbs  6,bitset ; Branch T bit was set
          ...
   bitset: nop           ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRCC - Branch if Carry Cleared

Description:

Conditional relative branch. Tests the Carry flag (C) and branches relatively to PC if C is cleared. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBC 0,k).

Operation:

(i) If C = 0 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRCC k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	01kk	kkkk	k000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            addr22,r23    ; Add r23 to r22
            brccnocarry   ; Branch if carry cleared
            ...
nocarry:    nop           ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRCS - Branch if Carry Set

Description:

Conditional relative branch. Tests the Carry flag (C) and branches relatively to PC if C is set. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBS 0,k).

Operation:

(i) If C = 1 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRCS k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	00kk	kkkk	k000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
          cpi   r26,$56   ; Compare r26 with $56
          brcs  carry     ; Branch if carry set
          ...
   carry: nop             ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BREQ - Branch if Equal

Description:

Conditional relative branch. Tests the Zero flag (Z) and branches relatively to PC if Z is set. If the instruction is executed immediately after any of the instructions CP, CPI, SUB or SUBI, the branch will occur if and only if the unsigned or signed binary number represented in Rd was equal to the unsigned or signed binary number represented in Rr. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBS 1,k).

Operation:

(i) If Rd = Rr (Z = 1) then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BREQ k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	00kk	kkkk	k001

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
          cpr1,r0    ; Compare registers r1 and r0
          breqequal  ; Branch if registers equal
          ...
   equal: nop        ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRGE - Branch if Greater or Equal (Signed)

Description:

Conditional relative branch. Tests the Signed flag (S) and branches relatively to PC if S is cleared. If the instruction is executed immediately after any of the instructions CP, CPI, SUB or SUBI, the branch will occur if and only if the signed binary number represented in Rd was greater than or equal to the signed binary number represented in Rr. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBC 4,k).

Operation:

(i) If Rd ≥ Rr (N ⊕ V = 0) then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRGE k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	01kk	kkkk	k100

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            cpr11,r12    ; Compare registers r11 and r12
            brgegreateq  ; Branch if r11 >= r12 (signed)
            ...
greateq:    nop          ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRHC - Branch if Half Carry Flag is Cleared

Description:

Conditional relative branch. Tests the Half Carry flag (H) and branches relatively to PC if H is cleared. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBC 5,k).

Operation:

(i) If H = 0 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRHC k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	01kk	kkkk	k101

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            brhc hclear      ; Branch if half carry flag cleared
            ...
hclear:     nop              ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRHS - Branch if Half Carry Flag is Set

Description:

Conditional relative branch. Tests the Half Carry flag (H) and branches relatively to PC if H is set. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBS 5,k).

Operation:

(i) If H = 1 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRHS k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	00kk	kkkk	k101

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            brhshset     ; Branch if half carry flag set
            ...
   hset:    nop          ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRID - Branch if Global Interrupt is Disabled

Description:

Conditional relative branch. Tests the Global Interrupt flag (I) and branches relatively to PC if I is cleared. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBC 7,k).

Operation:

(i) If I = 0 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRID k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	01kk	kkkk	k111

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            brid intdis    ; Branch if interrupt disabled
            ...
intdis:     nop            ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRIE - Branch if Global Interrupt is Enabled

Description:

Conditional relative branch. Tests the Global Interrupt flag (I) and branches relatively to PC if I is set. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBS 7,k).

Operation:

(i) If I = 1 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRIE k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	00kk	kkkk	k111

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            brieinten      ; Branch if interrupt enabled
            ...
inten:      nop            ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRLO - Branch if Lower (Unsigned)

Description:

Conditional relative branch. Tests the Carry flag (C) and branches relatively to PC if C is set. If the instruction is executed immediately after any of the instructions CP, CPI, SUB or SUBI, the branch will occur if and only if the unsigned binary number represented in Rd was smaller than the unsigned binary number represented in Rr. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBS 0,k).

Operation:

(i) If Rd < Rr (C = 1) then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRLO k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	00kk	kkkk	k000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
          eor   r19,r19    ; Clear r19
   loop:  inc   r19        ; Increase r19
          ...
          cpi   r19,$10    ; Compare r19 with $10
          brlo  loop       ; Branch if r19 < $10 (unsigned)
          nop              ; Exit from loop (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRLT - Branch if Less Than (Signed)

Description:

Conditional relative branch. Tests the Signed flag (S) and branches relatively to PC if S is set. If the instruction is executed immediately after any of the instructions CP, CPI, SUB or SUBI, the branch will occur if and only if the signed binary number represented in Rd was less than the signed binary number represented in Rr. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBS 4,k).

Operation:

(i) If Rd < Rr (N ⊕ V = 1) then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRLT k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	00kk	kkkk	k100

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            cp    r16,r1       ; Compare r16 to r1
            brlt  less         ; Branch if r16 < r1 (signed)
            ...
   less:    nop                ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRMI - Branch if Minus

Description:

Conditional relative branch. Tests the Negative flag (N) and branches relatively to PC if N is set. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBS 2,k).

Operation:

(i) If N = 1 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRMI k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	00kk	kkkk	k010

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            subi    r18,4       ; Subtract 4 from r18
            brmi    negative    ; Branch if result negative
            ...
negative:   nop                 ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRNE - Branch if Not Equal

Description:

Conditional relative branch. Tests the Zero flag (Z) and branches relatively to PC if Z is cleared. If the instruction is executed immediately after any of the instructions CP, CPI, SUB or SUBI, the branch will occur if and only if the unsigned or signed binary number represented in Rd was not equal to the unsigned or signed binary number represented in Rr. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBC 1,k).

Operation:

(i) If Rd ≠ Rr (Z = 0) then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRNE k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	01kk	kkkk	k001

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            eor     r27,r27   ; Clear r27
    loop:   inc     r27       ; Increase r27
            ...
            cpi     r27,5     ; Compare r27 to 5
            brne    loop      ; Branch if r27<>5
            nop               ; Loop exit (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRPL - Branch if Plus

Description:

Conditional relative branch. Tests the Negative flag (N) and branches relatively to PC if N is cleared. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBC 2,k).

Operation:

(i) If N = 0 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRPL k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	01kk	kkkk	k010

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
                subi  r26,$50        ; Subtract $50 from r26
                brpl  positive       ; Branch if r26 positive
                ...
   positive:    nop                  ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRSH - Branch if Same or Higher (Unsigned)

Description:

Conditional relative branch. Tests the Carry flag (C) and branches relatively to PC if C is cleared. If the instruction is executed immediately after execution of any of the instructions CP, CPI, SUB or SUBI the branch will occur if and only if the unsigned binary number represented in Rd was greater than or equal to the unsigned binary number represented in Rr. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBC 0,k).

Operation:

(i) If Rd ≥Rr (C = 0) then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRSH k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	01kk	kkkk	k000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            subi  r19,4        ; Subtract 4 from r19
            brsh  highsm       ; Branch if r19 >= 4 (unsigned)
            ...
highsm:     nop                ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRTC - Branch if the T Flag is Cleared

Description:

Conditional relative branch. Tests the T flag and branches relatively to PC if T is cleared. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBC 6,k).

Operation:

(i) If T = 0 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRTC k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	01kk	kkkk	k110

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            bst     r3,5      ; Store bit 5 of r3 in T flag
            brtc    tclear    ; Branch if this bit was cleared
            ...
tclear:     nop               ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRTS - Branch if the T Flag is Set

Description:

Conditional relative branch. Tests the T flag and branches relatively to PC if T is set. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBS 6,k).

Operation:

(i) If T = 1 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRTS k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	00kk	kkkk	k110

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            bst   r3,5     ; Store bit 5 of r3 in T flag
            brts  tset     ; Branch if this bit was set
            ...
   tset:    nop            ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BRVC - Branch if Overflow Cleared

Description:

Conditional relative branch. Tests the Overflow flag (V) and branches relatively to PC if V is cleared. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBC 3,k).

Operation:

(i) If V = 0 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRVC k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	01kk	kkkk	k011

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            add   r3,r4      ; Add r4 to r3
            brvc  noover     ; Branch if no overflow
            ...
noover:     nop              ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1 if condition is false
2 if condition is true

BRVS - Branch if Overflow Set

Description:

Conditional relative branch. Tests the Overflow flag (V) and branches relatively to PC if V is set. This instruction branches relatively to PC in either direction (PC-64≤destination≤PC+63). The parameter k is the offset from PC and is represented in two's complement form. (Equivalent to instruction BRBS 3,k).

Operation:

(i) If V = 1 then PC ← PC + k + 1, else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	BRVS k	-64 ≤ k ≤ +63	PC ← PC + k + 1 PC ← PC + 1, if condition is false

16 bit Opcode:

1111	00kk	kkkk	k011

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            add     r3,r4       ; Add r4 to r3
            brvs    overfl      ; Branch if overflow
            ...
overfl:     nop                 ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false
2 if condition is true

BSET - Bit Set in SREG

Description:
Sets a single flag or bit in SREG.

Operation:
(i) SREG(s) ← 1

	Syntax:	Operands:	Program Counter:
(i)	BSET s	0 ≤ s ≤ 7	PC ← PC + 1

16 bit Opcode:

1001	0100	0sss	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
⇔	⇔	⇔	⇔	⇔	⇔	⇔	⇔

I: 1 if s = 7; Unchanged otherwise.

T: 1 if s = 6; Unchanged otherwise.

H: 1 if s = 5; Unchanged otherwise.

S: 1 if s = 4; Unchanged otherwise.

V: 1 if s = 3; Unchanged otherwise.

N: 1 if s = 2; Unchanged otherwise.

Z: 1 if s = 1; Unchanged otherwise.

C: 1 if s = 0; Unchanged otherwise.

Example:
```
bset    6       ; Set T flag
bset    7       ; Enable interrupt
```

Words: 1 (2 bytes)
Cycles: 1

BST - Bit Store from Bit in Register to T Flag in SREG

Description:
Stores bit b from Rd to the T flag in SREG (status register).

Operation:
(i) T ← Rd(b)

	Syntax:	Operands:	Program Counter:
(i)	BST Rd,b	0 ≤ d ≤ 31, 0 ≤ b ≤ 7	PC ← PC + 1

16 bit Opcode:

1111	101d	dddd	Xbbb

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	⇔	-	-	-	-	-	-

T: 0 if bit b in Rd is cleared. Set to 1 otherwise.

Example:

```
                    ; Copy bit
bst     r1,2        ; Store bit 2 of r1 in T flag
bld     r0,4        ; Load T into bit 4 of r0
```

Words: 1 (2 bytes)
Cycles: 1

CALL - Long Call to a Subroutine

Description:

Calls to a subroutine within the entire program memory. The return address (to the instruction after the CALL) will be stored onto the stack. (See also RCALL).

	Operation:	
(i)	PC ← k	Devices with 16 bits PC, 128K bytes program memory maximum.
(ii)	PC ← k	Devices with 22 bits PC, 8M bytes program memory maximum.

	Syntax:	**Operands:**	**Program Counter:Stack**
(i)	CALL k	0 ≤ k ≤ 64K	PC ← kSTACK ← PC+2 SP ← SP-2, (2 bytes, 16 bits)
(ii)	CALL k	0 ≤ k ≤ 4M	PC ← kSTACK ← PC+2 SP ← SP-3 (3 bytes, 22 bits)

32 bit Opcode:

1001	010k	kkkk	111k
kkkk	kkkk	kkkk	kkkk

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
          mov    r16,r0       ; Copy r0 to r16
          call   check        ; Call subroutine
          nop                 ; Continue (do nothing)
          ...
check:    cpi    r16,$42      ; Check if r16 has a special value
          breq   error        ; Branch if equal
          ret                 ; Return from subroutine
          ...
error:    rjmp   error        ; Infinite loop
```

Words: 2 (4 bytes)
Cycles: 4

CBI - Clear Bit in I/O Register

Description:
Clears a specified bit in an I/O register. This instruction operates on the lower 32 I/O registers - addresses 0-31.

Operation:
(i) I/O(P,b) ← 0

	Syntax:	Operands:	Program Counter:
(i)	CBI P,b	0 ≤ P ≤ 31, 0 ≤ b ≤ 7	PC ← PC + 1

16 bit Opcode:

1001	1000	pppp	pbbb

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:
```
cbi     $12,7          ; Clear bit 7 in Port D
```

Words: 1 (2 bytes)
Cycles: 2

CBR - Clear Bits in Register

Description:

Clears the specified bits in register Rd. Performs the logical AND between the contents of register Rd and the complement of the constant mask K. The result will be placed in register Rd.

Operation:

(i) Rd ← Rd • ($FF - K)

	Syntax:	Operands:	Program Counter:
(i)	CBR Rd,K	16 ≤ d ≤ 31, 0 ≤ K ≤ 255	PC ← PC + 1

16 bit Opcode: See ANDI with K complemented.

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	**0**	⇔	⇔	-

S: N ⊕ V, For signed tests.

V: 0
Cleared

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
cbr     r16,$F0     ; Clear upper nibble of r16
cbr     r18,1       ; Clear bit 0 in r18
```

Words: 1 (2 bytes)
Cycles: 1

CLC - Clear Carry Flag

Description:
Clears the Carry flag (C) in SREG (status register).

Operation:
(i) C ← 0

	Syntax:	Operands:	Program Counter:
(i)	CLC	None	PC ← PC + 1

16 bit Opcode:

1001	0100	1000	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	**0**

C: 0
Carry flag cleared

Example:
```
add   r0,r0     ; Add r0 to itself
clc             ; Clear carry flag
```

Words: 1 (2 bytes)
Cycles: 1

CLH - Clear Half Carry Flag

Description:
Clears the Half Carry flag (H) in SREG (status register).

Operation:
(i) H ← 0

	Syntax:	Operands:	Program Counter:
(i)	CLH	None	PC ← PC + 1

16 bit Opcode:

1001	0100	1101	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	**0**	-	-	-	-	-

H: 0
Half Carry flag cleared

Example:
```
clh          ; Clear the Half Carry flag
```

Words: 1 (2 bytes)
Cycles: 1

CLI - Clear Global Interrupt Flag

Description:
Clears the Global Interrupt flag (I) in SREG (status register).

Operation:
(i) I ← 0

	Syntax:	Operands:	Program Counter:
(i)	CLI	None	PC ← PC + 1

16 bit Opcode:

1001	0100	1111	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
0	-	-	-	-	-	-	-

I: 0
Global Interrupt flag cleared

Example:

```
cli               ; Disable interrupts
in      r11,$16   ; Read port B
sei               ; Enable interrupts
```

Words: 1 (2 bytes)
Cycles: 1

CLN - Clear Negative Flag

Description:
Clears the Negative flag (N) in SREG (status register).

Operation:
(i) N ← 0

	Syntax:	Operands:	Program Counter:
(i)	CLN	None	PC ← PC + 1

16 bit Opcode:

1001	0100	1010	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	**0**	-	-

N: 0
Negative flag cleared

Example:
```
add     r2,r3     ; Add r3 to r2
cln               ; Clear negative flag
```

Words: 1 (2 bytes)
Cycles: 1

CLR - Clear Register

Description:

Clears a register. This instruction performs an Exclusive OR between a register and itself. This will clear all bits in the register.

Operation:

(i) Rd ← Rd ⊕ Rd

	Syntax:	Operands:	Program Counter:
(i)	CLR Rd	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode: (see EOR Rd,Rd)

0010	01dd	dddd	dddd

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	0	0	0	1	-

S: 0
Cleared

V: 0
Cleared

N: 0
Cleared

Z: 1
Set

R (Result) equals Rd after the operation.

Example:

```
          clr   r18        ; clear r18
   loop:  inc   r18        ; increase r18
          ...
          cpi   r18,$50    ; Compare r18 to $50
          brne  loop
```

Words: 1 (2 bytes)
Cycles: 1

CLS - Clear Signed Flag

Description:
Clears the Signed flag (S) in SREG (status register).

Operation:
(i) S ← 0

	Syntax:	Operands:	Program Counter:
(i)	CLS	None	PC ← PC + 1

16 bit Opcode:

1001	0100	1100	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	**0**	-	-	-	-

S: 0
Signed flag cleared

Example:
```
add   r2,r3      ; Add r3 to r2
cls              ; Clear signed flag
```

Words: 1 (2 bytes)
Cycles: 1

CLT - Clear T Flag

Description:
Clears the T flag in SREG (status register).

Operation:
(i) T ← 0

	Syntax:	Operands:	Program Counter:
(i)	CLT	None	PC ← PC + 1

16 bit Opcode:

1001	0100	1110	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	**0**	-	-	-	-	-	-

T: 0
T flag cleared

Example:
```
clt             ; Clear T flag
```

Words: 1 (2 bytes)
Cycles: 1

CLV - Clear Overflow Flag

Description:
Clears the Overflow flag (V) in SREG (status register).

Operation:
(i) V ← 0

	Syntax:	Operands:	Program Counter:
(i)	CLV	None	PC ← PC + 1

16 bit Opcode:

1001	0100	1011	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	**0**	-	-	-

V: 0
Overflow flag cleared

Example:

```
add     r2,r3     ; Add r3 to r2
clv               ; Clear overflow flag
```

Words: 1 (2 bytes)
Cycles: 1

CLZ - Clear Zero Flag

Description:
Clears the Zero flag (Z) in SREG (status register).

Operation:
(i) Z ← 0

	Syntax:	Operands:	Program Counter:
(i)	CLZ	None	PC ← PC + 1

16 bit Opcode:

1001	0100	1001	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	**0**	-

Z: 0
Zero flag cleared

Example:
```
add     r2,r3     ; Add r3 to r2
clz               ; Clear zero
```

Words: 1 (2 bytes)
Cycles: 1

COM - One's Complement

Description:
This instruction performs a one's complement of register Rd.

Operation:
(i) Rd ← $FF - Rd

	Syntax:	Operands:	Program Counter:
(i)	COM Rd	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode:

1001	010d	dddd	0000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	**0**	⇔	⇔	**1**

S: N ⊕ V
For signed tests.

V: 0
Cleared.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; Cleared otherwise.

C: 1
Set.

R (Result) equals Rd after the operation.

Example:
```
            com     r4          ; Take one's complement of r4
            breq    zero        ; Branch if zero
            ...
   zero:    nop                 ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

CP - Compare

Description:

This instruction performs a compare between two registers Rd and Rr. None of the registers are changed. All conditional branches can be used after this instruction.

Operation:

(i) Rd - Rr

	Syntax:	Operands:	Program Counter:
(i)	CP Rd,Rr	$0 \le d \le 31$, $0 \le r \le 31$	PC ← PC + 1

16 bit Opcode:

0001	01rd	dddd	rrrr

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: $\overline{Rd3} \bullet Rr3 + Rr3 \bullet R3 + R3 \bullet \overline{Rd3}$
Set if there was a borrow from bit 3; cleared otherwise

S: $N \oplus V$, For signed tests.

V: $Rd7 \bullet \overline{Rd7} \bullet \overline{R7} + \overline{Rd7} \bullet Rr7 \bullet R7$
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $Rd7 \bullet Rr7 + Rr7 \bullet R7 + R7 \bullet Rd7$
Set if the result is $00; cleared otherwise.

C: $\overline{Rd7} \bullet Rr7 + Rr7 \bullet R7 + R7 \bullet \overline{Rd7}$
Set if the absolute value of the contents of Rr is larger than the absolute value of Rd; cleared otherwise.

R (Result) after the operation.

Example:

```
          cp    r4,r19      ; Compare r4 with r19
          brne  noteq       ; Branch if r4 <> r19
          ...
   noteq: nop               ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

CPC - Compare with Carry

Description:

This instruction performs a compare between two registers Rd and Rr and also takes into account the previous carry. None of the registers are changed. All conditional branches can be used after this instruction.

Operation:

(i) Rd - Rr - C

	Syntax:	Operands:	Program Counter:
(i)	CPC Rd,Rr	$0 \le d \le 31, 0 \le r \le 31$	PC ← PC + 1

16 bit Opcode:

0000	01rd	dddd	rrrr

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: $\overline{Rd3} \bullet Rr3 + Rr3 \bullet R3 + R3 \bullet \overline{Rd3}$
Set if there was a borrow from bit 3; cleared otherwise

S: N ⊕ V, For signed tests.

V: $Rd7 \bullet \overline{Rr7} \bullet \overline{R7} + \overline{Rd7} \bullet Rr7 \bullet R7$
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0} \bullet Z$
Previous value remains unchanged when the result is zero; cleared otherwise.

C: $\overline{Rd7} \bullet Rr7 + Rr7 \bullet R7 + R7 \bullet \overline{Rd7}$
Set if the absolute value of the contents of Rr plus previous carry is larger than the absolute value of Rd; cleared otherwise.

R (Result) after the operation.

Example:

```
                                ; Compare r3:r2 with r1:r0
        cp      r2,r0           ; Compare low byte
        cpc     r3,r1           ; Compare high byte
        brne    noteq           ; Branch if not equal
        ...
noteq:  nop                     ; Branch destination (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1

CPI - Compare with Immediate

Description:

This instruction performs a compare between register Rd and a constant. The register is not changed. All conditional branches can be used after this instruction.

Operation:

(i) Rd - K

	Syntax:	Operands:	Program Counter:
(i)	CPI Rd,K	$16 \le d \le 31, 0 \le K \le 255$	PC ← PC + 1

16 bit Opcode:

0011	KKKK	dddd	KKKK

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: $\overline{Rd3} \bullet K3 + K3 \bullet R3 + R3 \bullet \overline{Rd3}$
Set if there was a borrow from bit 3; cleared otherwise

S: $N \oplus V$, For signed tests.

V: $Rd7 \bullet \overline{K7} \bullet \overline{R7} + \overline{Rd7} \bullet K7 \bullet R7$
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; cleared otherwise.

C: $\overline{Rd7} \bullet K7 + K7 \bullet R7 + R7 \bullet \overline{Rd7}$
Set if the absolute value of K is larger than the absolute value of Rd; cleared otherwise.

R (Result) after the operation.

Example:

```
            cpi     r19,3     ; Compare r19 with 3
            brne    error     ; Branch if r19<>3
            ...
   error:   nop               ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

CPSE - Compare Skip if Equal

Description:

This instruction performs a compare between two registers Rd and Rr, and skips the next instruction if Rd = Rr.

Operation:

(i) If Rd = Rr then PC ← PC + 2 (or 3) else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	CPSE Rd,Rr	0 ≤ d ≤ 31, 0 ≤ r ≤ 31	PC ← PC + 1, Condition false - no skip PC ← PC + 2, Skip a one word instruction PC ← PC + 3, Skip a two word instruction

16 bit Opcode:

0001	00rd	dddd	rrrr

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
inc     r4          ; Increase r4
cpse    r4,r0       ; Compare r4 to r0
neg     r4          ; Only executed if r4<>r0
nop                 ; Continue (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

DEC - Decrement

Description:

Subtracts one -1- from the contents of register Rd and places the result in the destination register Rd.

The C flag in SREG is not affected by the operation, thus allowing the DEC instruction to be used on a loop counter in multiple-precision computations.

When operating on unsigned values, only BREQ and BRNE branches can be expected to perform consistently. When operating on two's complement values, all signed branches are available.

Operation:

(i) Rd ← Rd - 1

	Syntax:	Operands:	Program Counter:
(i)	DEC Rd	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode:

1001	010d	dddd	1010

Status Register and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	⇔	⇔	⇔	-

S: N ⊕ V
For signed tests.

V: $\overline{R7}$ •R6 •R5 •R4• R3• R2 •R1• R0
Set if two's complement overflow resulted from the operation; cleared otherwise. Two's complement overflow occurs if and only if Rd was $80 before the operation.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}$ •$\overline{R6}$• $\overline{R5}$ •$\overline{R4}$• $\overline{R3}$• $\overline{R2}$• $\overline{R1}$• $\overline{R0}$
Set if the result is $00; Cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
          ldi   r17,$10      ; Load constant in r17
   loop:  add   r1,r2        ; Add r2 to r1
          dec   r17          ; Decrement r17
          brne  loop         ; Branch if r17<>0
          nop                ; Continue (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

EOR - Exclusive OR

Description:

Performs the logical EOR between the contents of register Rd and register Rr and places the result in the destination register Rd.

Operation:

(i) Rd ← Rd ⊕ Rr

	Syntax:	Operands:	Program Counter:
(i)	EOR Rd,Rr	0 ≤ d ≤ 31, 0 ≤ r ≤ 31	PC ← PC + 1

16 bit Opcode:

0010	01rd	dddd	rrrr

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	**0**	⇔	⇔	-

S: N ⊕ V, For signed tests.

V: 0
Cleared

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
eor     r4,r4          ; Clear r4
eor     r0,r22         ; Bitwise exclusive or between r0 and r22
```

Words: 1 (2 bytes)
Cycles: 1

ICALL - Indirect Call to Subroutine

Description:

Indirect call of a subroutine pointed to by the Z (16 bits) pointer register in the register file. The Z pointer register is 16 bits wide and allows call to a subroutine within the current 64K words (128K bytes) section in the program memory space.

Operation:

(i) PC(15-0) ← Z(15 - 0)Devices with 16 bits PC, 128K bytes program memory maximum.

(ii) PC(15-0) ← Z(15 - 0)Devices with 22 bits PC, 8M bytes program memory maximum.
PC(21-16) is unchanged

	Syntax:	Operands:	Program Counter:	Stack
(i)	ICALL	None	See Operation	STACK ← PC+1 SP ← SP-2 (2 bytes, 16 bits)
(ii)	ICALL	None	See Operation	STACK ← PC+1 SP ← SP-3 (3 bytes, 22 bits)

16 bit Opcode:

1001	0101	XXXX	1001

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
mov     r30,r0      ; Set offset to call table
icall               ; Call routine pointed to by r31:r30
```

Words: 1 (2 bytes)

Cycles: 3

IJMP - Indirect Jump

Description:

Indirect jump to the address pointed to by the Z (16 bits) pointer register in the register file. The Z pointer register is 16 bits wide and allows jump within the current 64K words (128K bytes) section of program memory.

Operation:

(i) PC ← Z(15 - 0) Devices with 16 bits PC, 128K bytes program memory maximum.

(ii) PC(15-0) ← Z(15-0) Devices with 22 bits PC, 8M bytes program memory maximum.
PC(21-16) is unchanged

	Syntax:	Operands:	Program Counter:	Stack
(ii)	IJMP	None	See Operation	Not Affected
(iii)	IJMP	None	See Operation	Not Affected

16 bit Opcode:

1001	0100	XXXX	1001

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
mov     r30,r0    ; Set offset to jump table
ijmp              ; Jump to routine pointed to by r31:r30
```

Words: 1 (2 bytes)
Cycles: 2

IN - Load an I/O Port to Register

Description:
Loads data from the I/O Space (Ports, Timers, Configuration registers etc.) into register Rd in the register file.

Operation:
(i) Rd ← P

	Syntax:	Operands:	Program Counter:
(i)	IN Rd,P	0 ≤ d ≤ 31, 0 ≤ P ≤ 63	PC ← PC + 1

16 bit Opcode:

1011	0PPd	dddd	PPPP

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            in      r25,$16     ; Read Port B
            cpi     r25,4       ; Compare read value to constant
            breq    exit        ; Branch if r25=4
            ...
    exit:   nop                 ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

INC - Increment

Description:

Adds one -1- to the contents of register Rd and places the result in the destination register Rd.

The C flag in SREG is not affected by the operation, thus allowing the INC instruction to be used on a loop counter in multiple-precision computations.

When operating on unsigned numbers, only BREQ and BRNE branches can be expected to perform consistently. When operating on two's complement values, all signed branches are available.

Operation:

(i) Rd ← Rd + 1

	Syntax:	Operands:	Program Counter:
(i)	INC Rd	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode:

1001	010d	dddd	0011

Status Register and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	⇔	⇔	⇔	-

S: N ⊕ V

For signed tests.

V: R7 •$\overline{R6}$ •$\overline{R5}$ •$\overline{R4}$ •$\overline{R3}$• $\overline{R2}$ •$\overline{R1}$ •$\overline{R0}$
Set if two's complement overflow resulted from the operation; cleared otherwise. Two's complement overflow occurs if and only if Rd was $7F before the operation.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}$ •$\overline{R6}$ •$\overline{R5}$ •$\overline{R4}$•$\overline{R3}$ •$\overline{R2}$• $\overline{R1}$• $\overline{R0}$
Set if the result is $00; Cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
          clr     r22         ; clear r22
   loop:  inc     r22         ; increment r22
          ...
          cpi     r22,$4F     ; Compare r22 to $4f
          brne    loop        ; Branch if not equal
          nop                 ; Continue (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

JMP - Jump

Description:
Jump to an address within the entire 4M (words) program memory. See also RJMP.

Operation:
(i) PC ← k

	Syntax:	Operands:	Program Counter:	Stack
(i)	JMP k	0 ≤ k ≤ 4M	PC ← k	Unchanged

32 bit Opcode:

1001	010k	kkkk	110k
kkkk	kkkk	kkkk	kkkk

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
          mov     r1,r0     ; Copy r0 to r1
          jmp     farplc    ; Unconditional jump
          ...
farplc:   nop               ; Jump destination (do nothing)
```

Words: 2 (4 bytes)
Cycles: 3

LD - Load Indirect from SRAM to Register using Index X

Description:

Loads one byte indirect from SRAM to register. The SRAM location is pointed to by the X (16 bits) pointer register in the register file. Memory access is limited to the current SRAM page of 64K bytes. To access another SRAM page the RAMPX in register in the I/O area has to be changed.

The X pointer register can either be left unchanged after the operation, or it can be incremented or decremented. These features are especially suited for accessing arrays, tables, and stack pointer usage of the X pointer register.

Using the X pointer:

	Operation:	**Comment:**	
(i)	Rd ← (X)		X: Unchanged
(ii)	Rd ← (X)	X ← X + 1	X: Post incremented
(iii)	X ← X - 1	Rd ← (X)	X: Pre decremented

	Syntax:	**Operands:**	**Program Counter:**
(i)	LD Rd, X	0 ≤ d ≤ 31	PC ← PC + 1
(ii)	LD Rd, X+	0 ≤ d ≤ 31	PC ← PC + 1
(iii)	LD Rd,-X	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode :

(i)	1001	000d	dddd	1100
(ii)	1001	000d	dddd	1101
(iii)	1001	000d	dddd	1110

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
clr   r27           ; Clear X high byte
ldi   r26,$20       ; Set X low byte to $20
ld    r0,X+         ; Load r0 with SRAM loc. $20(X post inc)
ld    r1,X          ; Load r1 with SRAM loc. $21
ldi   r26,$23       ; Set X low byte to $23
ld    r2,X          ; Load r2 with SRAM loc. $23
ld    r3,-X         ; Load r3 with SRAM loc. $22(X pre dec)
```

Words: 1 (2 bytes)
Cycles: 2

LD (LDD) - Load Indirect from SRAM to Register using Index Y

Description:

Loads one byte indirect with or without displacement from SRAM to register. The SRAM location is pointed to by the Y (16 bits) pointer register in the register file. Memory access is limited to the current SRAM page of 64K bytes. To access another SRAM page the RAMPY register in the I/O area has to be changed.

The Y pointer register can either be left unchanged after the operation, or it can be incremented or decremented. These features are especially suited for accessing arrays, tables, and stack pointer usage of the Y pointer register.

Using the Y pointer:

	Operation:		**Comment:**
(i)	Rd ← (Y)		Y: Unchanged
(ii)	Rd ← (Y)	Y ← Y + 1	Y: Post incremented
(iii)	Y ← Y - 1	Rd ← (Y)	Y: Pre decremented
(iiii)	Rd ← (Y+q)		Y: Unchanged, q: Displacement

	Syntax:	**Operands:**	**Program Counter:**
(i)	LD Rd, Y	0 ≤ d ≤ 31	PC ← PC + 1
(ii)	LD Rd, Y+	0 ≤ d ≤ 31	PC ← PC + 1
(iii)	LD Rd,-Y	0 ≤ d ≤ 31	PC ← PC + 1
(iiii)	LDD Rd, Y+q	0 ≤ d ≤ 31, 0 ≤ q ≤ 63	PC ← PC + 1

16 bit Opcode :

(i)	1000	000d	dddd	1000
(ii)	1001	000d	dddd	1001
(iii)	1001	000d	dddd	1010
(iiii)	10q0	qq0d	dddd	1qqq

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
clr   r29          ; Clear Y high byte
ldi   r28,$20      ; Set Y low byte to $20
ld    r0,Y+        ; Load r0 with SRAM loc. $20(Y post inc)
ld    r1,Y         ; Load r1 with SRAM loc. $21
ldi   r28,$23      ; Set Y low byte to $23
ld    r2,Y         ; Load r2 with SRAM loc. $23
ld    r3,-Y        ; Load r3 with SRAM loc. $22(Y pre dec)
ldd   r4,Y+2       ; Load r4 with SRAM loc. $24
```

Words: 1 (2 bytes)
Cycles: 2

LD (LDD) - Load Indirect From SRAM to Register using Index Z

Description:

Loads one byte indirectly with or without displacement from SRAM to register. The SRAM location is pointed to by the Z (16 bits) pointer register in the register file. Memory access is limited to the current SRAM page of 64K bytes. To access another SRAM page the RAMPZ register in the I/O area has to be changed.

The Z pointer register can either be left unchanged after the operation, or it can be incremented or decremented. These features are especially suited for stack pointer usage of the Z pointer register, however because the Z pointer register can be used for indirect subroutine calls, indirect jumps and table lookup, it is often more convenient to use the X or Y pointer as a dedicated stack pointer.

For using the Z pointer for table lookup in program memory see the LPM instruction.

Using the Z pointer:

	Operation:	**Comment:**	
(i)	Rd ← (Z)		Z: Unchanged
(ii)	Rd ← (Z)	Z ← Z + 1	Z: Post increment
(iii)	Z ← Z -1	Rd ← (Z)	Z: Pre decrement
(iiii)	Rd ← (Z+q)		Z: Unchanged, q: Displacement

	Syntax:	**Operands:**	**Program Counter:**
(i)	LD Rd, Z	$0 \le d \le 31$	PC ← PC + 1
(ii)	LD Rd, Z+	$0 \le d \le 31$	PC ← PC + 1
(iii)	LD Rd,-Z	$0 \le d \le 31$	PC ← PC + 1
(iiii)	LDD Rd, Z+q	$0 \le d \le 31$, $0 \le q \le 63$	PC ← PC + 1

16 bit Opcode :

(i)	1000	000d	dddd	0000
(ii)	1001	000d	dddd	0001
(iii)	1001	000d	dddd	0010
(iiii)	10q0	qq0d	dddd	0qqq

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
clr   r31        ; Clear Z high byte
ldi   r30,$20    ; Set Z low byte to $20
ld    r0,Z+      ; Load r0 with SRAM loc. $20(Z post inc)
ld    r1,Z       ; Load r1 with SRAM loc. $21
ldi   r30,$23    ; Set Z low byte to $23
ld    r2,Z       ; Load r2 with SRAM loc. $23
ld    r3,-Z      ; Load r3 with SRAM loc. $22(Z pre dec)
ldd   r4,Z+2     ; Load r4 with SRAM loc. $24
```

Words: 1 (2 bytes)
Cycles: 2

LDI - Load Immediate

Description:
Loads an 8 bit constant directly to register 16 to 31.

Operation:
(i) Rd ← K

	Syntax:	Operands:	Program Counter:
(i)	LDI Rd,K	16 ≤ d ≤ 31, 0 ≤ K ≤ 255	PC ← PC + 1

16 bit Opcode:

1110	KKKK	dddd	KKKK

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
clr   r31          ; Clear Z high byte
ldi   r30,$F0      ; Set Z low byte to $F0
lpm                ; Load constant from program
                   ; memory pointed to by Z
```

Words: 1 (2 bytes)
Cycles: 1

LDS - Load Direct from SRAM

Description:

Loads one byte from the SRAM to a Register. A 16-bit address must be supplied. Memory access is limited to the current SRAM Page of 64K bytes. The LDS instruction uses the RAMPZ register to access memory above 64K bytes.

Operation:

(i) Rd ← (k)

	Syntax:	Operands:	Program Counter:
(i)	LDS Rd,k	0 ≤ d ≤ 31, 0 ≤ k ≤ 65535	PC ← PC + 2

32 bit Opcode:

1001	000d	dddd	0000
kkkk	kkkk	kkkk	kkkk

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
lds   r2,$FF00     ; Load r2 with the contents of SRAM location $FF00
add   r2,r1        ; add r1 to r2
sts   $FF00,r2     ; Write back
```

Words: 2 (4 bytes)
Cycles: 3

LPM - Load Program Memory

Description:

Loads one byte pointed to by the Z register into register 0 (R0). This instruction features a 100% space effective constant initialization or constant data fetch. The program memory is organized in 16 bits words and the LSB of the Z (16 bits) pointer selects either low byte (0) or high byte (1). This instruction can address the first 64K bytes (32K words) of program memory.

	Operation:	Comment:
(i)	R0 ← (Z)	Z points to program memory

	Syntax:	Operands:	Program Counter:
(i)	LPM	None	PC ← PC + 1

16 bit Opcode:

1001	0101	110X	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
clr   r31           ; Clear Z high byte
ldi   r30,$F0       ; Set Z low byte
lpm                 ; Load constant from program
                    ; memory pointed to by Z (r31:r30)
```

Words: 1 (2 bytes)
Cycles: 3

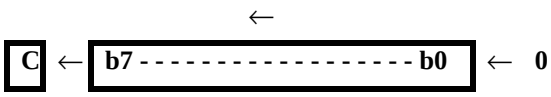
LSL - Logical Shift Left

Description:

Shifts all bits in Rd one place to the left. Bit 0 is cleared. Bit 7 is loaded into the C flag of the SREG. This operation effectively multiplies an unsigned value by two.

Operation:

(i)



	Syntax:	Operands:	Program Counter:
(i)	LSL Rd	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode: (see ADD Rd,Rd)

0000	11dd	dddd	dddd

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: Rd3

S: N ⊕ V, For signed tests.

V: N ⊕ C (For N and C after the shift)
Set if (N is set and C is clear) or (N is clear and C is set); Cleared otherwise (for values of N and C after the shift).

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; cleared otherwise.

C: Rd7
Set if, before the shift, the MSB of Rd was set; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
add     r0,r4       ; Add r4 to r0
lsl     r0          ; Multiply r0 by 2
```

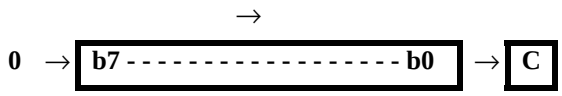
Words: 1 (2 bytes)
Cycles: 1

LSR - Logical Shift Right

Description:

Shifts all bits in Rd one place to the right. Bit 7 is cleared. Bit 0 is loaded into the C flag of the SREG. This operation effectively divides an unsigned value by two. The C flag can be used to round the result.

Operation:



	Syntax:	Operands:	Program Counter:
(i)	LSR Rd	$0 \le d \le 31$	PC ← PC + 1

16 bit Opcode:

1001	010d	dddd	0110

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	⇔	**0**	⇔	⇔

S: $N \oplus V$, For signed tests.

V: $N \oplus C$ (For N and C after the shift)
Set if (N is set and C is clear) or (N is clear and C is set); Cleared otherwise (for values of N and C after the shift).

N: 0

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; cleared otherwise.

C: Rd0
Set if, before the shift, the LSB of Rd was set; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
add     r0,r4          ; Add r4 to r0
lsr     r0             ; Divide r0 by 2
```

Words: 1 (2 bytes)
Cycles: 1

MOV - Copy Register

Description:

This instruction makes a copy of one register into another. The source register Rr is left unchanged, while the destination register Rd is loaded with a copy of Rr.

Operation:

(i) Rd ← Rr

	Syntax:	Operands:	Program Counter:
(i)	MOV Rd,Rr	0 ≤ d ≤ 31, 0 ≤ r ≤ 31	PC ← PC + 1

16 bit Opcode:

0010	11rd	dddd	rrrr

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

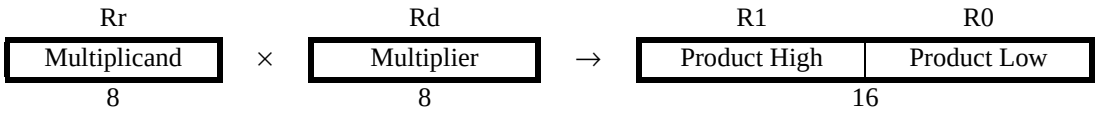
```
            mov     r16,r0    ; Copy r0 to r16
            call    check     ; Call subroutine
            ...
   check:   cpi     r16,$11   ; Compare r16 to $11
            ...
            ret               ; Return from subroutine
```

Words: 1 (2 bytes)
Cycles: 1

MUL - Multiply

Description:

This instruction performs 8-bit × 8-bit → 16-bit unsigned multiplication.



The multiplicand Rr and the multiplier Rd are two registers. The 16-bit product is placed in R1 (high byte) and R0 (low byte). Note that if the multiplicand and the multiplier is selected from R0 or R1 the result will overwrite those after multiplication.

Operation:

(i) R1,R0 ← Rr × Rd

	Syntax:	Operands:	Program Counter:
(i)	MUL Rd,Rr	0 ≤ d ≤ 31, 0 ≤ r ≤ 31	PC ← PC + 1

16 bit Opcode:

1001	11rd	dddd	rrrr

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	⇔

C: R15
Set if bit 15 of the result is set; cleared otherwise.

R (Result) equals R1,R0 after the operation.

Example:

```
mulr6,r5; Multiply r6 and r5
movr6,r1; Copy result back in r6:r5
movr5,r0; Copy result back in r6:r5
```

Words: 1 (2 bytes)
Cycles: 2

Not available in base-line microcontrollers.

NEG - Two's Complement

Description:

Replaces the contents of register Rd with its two's complement; the value $80 is left unchanged.

Operation:

(i) Rd ← $00 - Rd

	Syntax:	Operands:	Program Counter:
(i)	NEG Rd	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode:

1001	010d	dddd	0001

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: $R3 \bullet \overline{Rd3}$
Set if there was a borrow from bit 3; cleared otherwise

S: $N \oplus V$
For signed tests.

V: $R7 \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if there is a two's complement overflow from the implied subtraction from zero; cleared otherwise. A two's complement overflow will occur if and only if the contents of the Register after operation (Result) is $80.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; Cleared otherwise.

C: R7 + R6 + R5 + R4 + R3 + R2 + R1 + R0
Set if there is a borrow in the implied subtraction from zero; cleared otherwise. The C flag will be set in all cases except when the contents of Register after operation is $00.

R (Result) equals Rd after the operation.

Example:

```
            sub   r11,r0       ; Subtract r0 from r11
            brpl  positive     ; Branch if result positive
            neg   r11          ; Take two's complement of r11
positive:   nop                ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

NOP - No Operation

Description:
This instruction performs a single cycle No Operation.

Operation:
(i) No

	Syntax:	Operands:	Program Counter:
(i)	NOP	None	PC ← PC + 1

16 bit Opcode:

0000	0000	0000	0000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
clr     r16         ; Clear r16
ser     r17         ; Set r17
out     $18,r16     ; Write zeros to Port B
nop                 ; Wait (do nothing)
out     $18,r17     ; Write ones to Port B
```

Words: 1 (2 bytes)
Cycles: 1

OR - Logical OR

Description:

Performs the logical OR between the contents of register Rd and register Rr and places the result in the destination register Rd.

Operation:

(i) Rd ← Rd v Rr

	Syntax:	Operands:	Program Counter:
(i)	OR Rd,Rr	0 ≤ d ≤ 31, 0 ≤ r ≤ 31	PC ← PC + 1

16 bit Opcode:

0010	10rd	dddd	rrrr

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	**0**	⇔	⇔	-

S: N ⊕ V, For signed tests.

V: 0
Cleared

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}\bullet\ \overline{R6}\ \bullet\overline{R5}\bullet\ \overline{R4}\bullet\ \overline{R3}\ \bullet\overline{R2}\bullet\ \overline{R1}\bullet\ \overline{R0}$
Set if the result is $00; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
          or     r15,r16      ; Do bitwise or between registers
          bst    r15,6        ; Store bit 6 of r15 in T flag
          brts   ok           ; Branch if T flag set
          ...
   ok:    nop                 ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

ORI - Logical OR with Immediate

Description:

Performs the logical OR between the contents of register Rd and a constant and places the result in the destination register Rd.

Operation:

(i) Rd ← Rd v K

	Syntax:	Operands:	Program Counter:
(i)	ORI Rd,K	16 ≤ d ≤ 31, 0 ≤ K ≤ 255	PC ← PC + 1

16 bit Opcode:

0110	KKKK	dddd	KKKK

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	0	⇔	⇔	-

S: N ⊕ V, For signed tests.

V: 0
Cleared

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
ori     r16,$F0         ; Set high nibble of r16
ori     r17,1           ; Set bit 0 of r17
```

Words: 1 (2 bytes)
Cycles: 1

OUT - Store Register to I/O port

Description:
Stores data from register Rr in the register file to I/O space (Ports, Timers, Configuration registers etc.).

Operation:
(i) P ← Rr

	Syntax:	Operands:	Program Counter:
(i)	OUT P,Rr	0 ≤ r ≤ 31, 0 ≤ P ≤ 63	PC ← PC + 1

16 bit Opcode:

1011	1PPr	rrrr	PPPP

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:
```
clr   r16           ; Clear r16
ser   r17           ; Set r17
out   $18,r16       ; Write zeros to Port B
nop                 ; Wait (do nothing)
out   $18,r17       ; Write ones to Port B
```

Words: 1 (2 bytes)
Cycles: 1

POP - Pop Register from Stack

Description:
This instruction loads register Rd with a byte from the STACK.

Operation:
(i) Rd ← STACK

	Syntax:	Operands:	Program Counter:Stack
(i)	POP Rd	0 ≤ d ≤ 31	PC ← PC + 1SP ← SP + 1

16 bit Opcode:

1001	000d	dddd	1111

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
                call    routine    ; Call subroutine
                ...
routine:        push    r14        ; Save r14 on the stack
                push    r13        ; Save r13 on the stack
                ...
                pop     r13        ; Restore r13
                pop     r14        ; Restore r14
                ret                ; Return from subroutine
```

Words: 1 (2 bytes)
Cycles: 2

PUSH - Push Register on Stack

Description:
This instruction stores the contents of register Rr on the STACK.

Operation:
(i) STACK ← Rr

	Syntax:	Operands:	Program Counter:Stack:
(i)	PUSH Rr	0 ≤ r ≤ 31	PC ← PC + 1SP ← SP - 1

16 bit Opcode:

1001	001d	dddd	1111

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            call    routine  ; Call subroutine
            ...
routine:    push    r14      ; Save r14 on the stack
            push    r13      ; Save r13 on the stack
            ...
            pop     r13      ; Restore r13
            pop     r14      ; Restore r14
            ret              ; Return from subroutine
```

Words: 1 (2 bytes)
Cycles: 2

RCALL - Relative Call to Subroutine

Description:

Calls a subroutine within ± 2K words (4K bytes). The return address (the instruction after the RCALL) is stored onto the stack. (See also CALL).

Operation:

(i) PC ← PC + k + 1 Devices with 16 bits PC, 128K bytes program memory maximum.
(ii) PC ← PC + k + 1 Devices with 22 bits PC, 8M bytes program memory maximum.

	Syntax:	Operands:	Program Counter:	Stack
(i)	RCALL k	-2K ≤ k ≤ 2K	PC ← PC + k + 1	STACK ← PC+1 SP ← SP-2 (2 bytes, 16 bits)
(ii)	RCALL k	-2K ≤ k ≤ 2K	PC ← PC + k + 1	STACK ← PC+1 SP ← SP-3 (3 bytes, 22 bits)

16 bit Opcode:

1101	kkkk	kkkk	kkkk

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            rcall   routine      ; Call subroutine
            ...
routine:    push    r14          ; Save r14 on the stack
            ...
            pop     r14          ; Restore r14
            ret                  ; Return from subroutine
```

Words: 1 (2 bytes)
Cycles: 3

RET - Return from Subroutine

Description:

Returns from subroutine. The return address is loaded from the STACK.

Operation:

(i) PC(15-0) ← STACK Devices with 16 bits PC, 128K bytes program memory maximum.

(ii) PC(21-0) ← STACK Devices with 22 bits PC, 8M bytes program memory maximum.

	Syntax:	Operands:	Program Counter:	Stack
(i)	RET	None	See Operation	SP←SP+2,(2 bytes,16 bits pulled)
(ii)	RET	None	See Operation	SP←SP+3,(3 bytes,22 bits pulled)

16 bit Opcode:

1001	0101	0XX0	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
                call    routine     ; Call subroutine
                ...
routine:        push    r14         ; Save r14 on the stack
                ...
                pop     r14         ; Restore r14
                ret                 ; Return from subroutine
```

Words: 1 (2 bytes)

Cycles: 4

RETI - Return from Interrupt

Description:

Returns from interrupt. The return address is loaded from the STACK and the global interrupt flag is set.

Operation:

(i) PC(15-0) ← STACK Devices with 16 bits PC, 128K bytes program memory maximum.

(ii) PC(21-0) ← STACK Devices with 22 bits PC, 8M bytes program memory maximum.

	Syntax:	Operands:	Program Counter:	Stack
(i)	RETI	None	See Operation	SP ← SP +2 (2 bytes, 16 bits)
(ii)	RETI	None	See Operation	SP ← SP +3 (3 bytes, 22 bits)

16 bit Opcode:

1001	0101	0XX1	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
1	-	-	-	-	-	-	-

I: 1
The I flag is set.

Example:

```
            ...
extint:     push    r0        ; Save r0 on the stack
            ...
            pop     r0        ; Restore r0
            reti              ; Return and enable interrupts
```

Words: 1 (2 bytes)
Cycles: 4

RJMP - Relative Jump

Description:

Relative jump to an address within PC-2K and PC + 2K (words). In the assembler, labels are used instead of relative operands. For AVR microcontrollers with program memory not exceeding 4K words (8K bytes) this instruction can address the entire memory from every address location.

Operation:

(i) PC ← PC + k + 1

	Syntax:	Operands:	Program Counter:	Stack
(i)	RJMP k	-2K ≤ k ≤ 2K	PC ← PC + k + 1	Unchanged

16 bit Opcode:

1100	kkkk	kkkk	kkkk

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
            cpi     r16,$42      ; Compare r16 to $42
            brne    error        ; Branch if r16 <> $42
            rjmp    ok           ; Unconditional branch
    error:  add     r16,r17      ; Add r17 to r16
            inc     r16          ; Increment r16
    ok:     nop                  ; Destination for rjmp (do nothing)
```

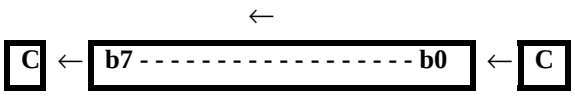
Words: 1 (2 bytes)
Cycles: 2

ROL - Rotate Left trough Carry

Description:

Shifts all bits in Rd one place to the left. The C flag is shifted into bit 0 of Rd. Bit 7 is shifted into the C flag.

Operation:



	Syntax:	Operands:	Program Counter:
(i)	ROL Rd	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode: (see ADC Rd,Rd)

0001	11dd	dddd	dddd

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: Rd3

S: N ⊕ V, For signed tests.

V: N ⊕ C (For N and C after the shift)
Set if (N is set and C is clear) or (N is clear and C is set); Cleared otherwise (for values of N and C after the shift).

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}$• $\overline{R6}$ •$\overline{R5}$• $\overline{R4}$• $\overline{R3}$ •$\overline{R2}$• $\overline{R1}$• $\overline{R0}$
Set if the result is $00; cleared otherwise.

C: Rd7
Set if, before the shift, the MSB of Rd was set; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
            rolr15          ; Rotate left
            brcsoneenc      ; Branch if carry set
            ...
oneenc:     nop             ; Branch destination (do nothing)
```

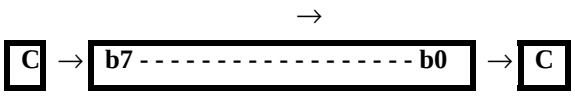
Words: 1 (2 bytes)
Cycles: 1

ROR - Rotate Right trough Carry

Description:

Shifts all bits in Rd one place to the right. The C flag is shifted into bit 7 of Rd. Bit 0 is shifted into the C flag.

Operation:



	Syntax:	Operands:	Program Counter:
(i)	ROR Rd	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode:

1001	010d	dddd	0111

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	⇔	⇔	⇔	⇔

S: N ⊕ V, For signed tests.

V: N ⊕ C (For N and C after the shift)
Set if (N is set and C is clear) or (N is clear and C is set); Cleared otherwise (for values of N and C after the shift).

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}$• $\overline{R6}$ •$\overline{R5}$• $\overline{R4}$• $\overline{R3}$ •$\overline{R2}$• $\overline{R1}$• $\overline{R0}$
Set if the result is $00; cleared otherwise.

C: Rd0
Set if, before the shift, the LSB of Rd was set; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
            rorr15          ; Rotate right
            brcczeroenc     ; Branch if carry cleared
            ...
zeroenc:    nop             ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

SBC - Subtract with Carry

Description:

Subtracts two registers and subtracts with the C flag and places the result in the destination register Rd.

Operation:

(i) Rd ← Rd - Rr - C

	Syntax:	Operands:	Program Counter:
(i)	SBC Rd,Rr	$0 \le d \le 31, 0 \le r \le 31$	PC ← PC + 1

16 bit Opcode:

0000	10rd	dddd	rrrr

Status Register and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: $\overline{Rd3}$• Rr3 + Rr3• R3 + R3 •$\overline{Rd3}$
Set if there was a borrow from bit 3; cleared otherwise

S: N ⊕ V, For signed tests.

V: Rd7 •$\overline{Rr7}$• $\overline{R7}$ +$\overline{Rd7}$ •Rr7 •R7
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}$• $\overline{R6}$ •$\overline{R5}$• $\overline{R4}$• $\overline{R3}$ •$\overline{R2}$• $\overline{R1}$• $\overline{R0}$• Z
Previous value remains unchanged when the result is zero; cleared otherwise.

C: $\overline{Rd7}$ •Rr7+ Rr7 •R7 +R7 •$\overline{Rd7}$
Set if the absolute value of the contents of Rr plus previous carry is larger than the absolute value of the Rd; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
                              ; Subtract r1:r0 from r3:r2
        sub     r2,r0         ; Subtract low byte
        sbc     r3,r1         ; Subtract with carry high byte
```

Words: 1 (2 bytes)
Cycles: 1

SBCI - Subtract Immediate with Carry

Description:

Subtracts a constant from a register and subtracts with the C flag and places the result in the destination register Rd.

Operation:

(i) Rd ← Rd - K - C

	Syntax:	Operands:	Program Counter:
(i)	SBCI Rd,K	16 ≤ d ≤ 31, 0 ≤ K ≤ 255	PC ← PC + 1

16 bit Opcode:

0100	KKKK	dddd	KKKK

Status Register and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: $\overline{Rd3}$• K3 + K3• R3 + R3 •$\overline{Rd3}$
Set if there was a borrow from bit 3; cleared otherwise

S: N ⊕ V, For signed tests.

V: Rd7 •$\overline{K7}$• $\overline{R7}$ +$\overline{Rd7}$ •K7 •R7
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}$• $\overline{R6}$ •$\overline{R5}$• $\overline{R4}$• $\overline{R3}$ •$\overline{R2}$• $\overline{R1}$• $\overline{R0}$• Z
Previous value remains unchanged when the result is zero; cleared otherwise.

C: $\overline{Rd7}$ •K7+ K7 • R7 +R7 •$\overline{Rd7}$
Set if the absolute value of the constant plus previous carry is larger than the absolute value of Rd; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
                        ; Subtract $4F23 from r17:r16
        subi  r16,$23   ; Subtract low byte
        sbci  r17,$4F   ; Subtract with carry high byte
```

Words: 1 (2 bytes)
Cycles: 1

SBI - Set Bit in I/O Register

Description:
Sets a specified bit in an I/O register. This instruction operates on the lower 32 I/O registers - addresses 0-31.

Operation:
(i) I/O(P,b) ← 1

	Syntax:	Operands:	Program Counter:
(i)	SBI P,b	0 ≤ P ≤ 31, 0 ≤ b ≤ 7	PC ← PC + 1

16 bit Opcode:

1001	1010	pppp	pbbb

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:
```
out   $1E,r0        ; Write EEPROM address
sbi   $1C,0         ; Set read bit in EECR
in    r1,$1D        ; Read EEPROM data
```

Words: 1 (2 bytes)
Cycles: 2

SBIC - Skip if Bit in I/O Register is Cleared

Description:

This instruction tests a single bit in an I/O register and skips the next instruction if the bit is cleared. This instruction operates on the lower 32 I/O registers - addresses 0-31.

Operation:

(i) If I/O(P,b) = 0 then PC ← PC + 2 (or 3) else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	SBIC P,b	0 ≤ P ≤ 31, 0 ≤ b ≤ 7	PC ← PC + 1, If condition is false, no skip. PC ← PC + 2, If next instruction is one word. PC ← PC + 3, If next instruction is JMP or CALL

16 bit Opcode:

1001	1001	pppp	pbbb

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
e2wait:  sbic  $1C,1      ; Skip next inst. if EEWE cleared
         rjmp  e2wait     ; EEPROM write not finished
         nop              ; Continue (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false (no skip)
2 if condition is true (skip is executed)

SBIS - Skip if Bit in I/O Register is Set

Description:

This instruction tests a single bit in an I/O register and skips the next instruction if the bit is set. This instruction operates on the lower 32 I/O registers - addresses 0-31.

Operation:

(i) If I/O(P,b) = 1 then PC ← PC + 2 (or 3) else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	SBIS P,b	0 ≤ P ≤ 31, 0 ≤ b ≤ 7	PC ← PC + 1, Condition false - no skip PC ← PC + 2, Skip a one word instruction PC ← PC + 3, Skip a JMP or a CALL

16 bit Opcode:

1001	1011	pppp	pbbb

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
waitset: sbis  $10,0          ; Skip next inst. if bit 0 in Port D set
         rjmp  waitset        ; Bit not set
         nop                  ; Continue (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false (no skip)
2 if condition is true (skip is executed)

SBIW - Subtract Immediate from Word

Description:

Subtracts an immediate value (0-63) from a register pair and places the result in the register pair. This instruction operates on the upper four register pairs, and is well suited for operations on the pointer registers.

Operation:

(i) Rdh:Rdl ← Rdh:Rdl - K

	Syntax:	Operands:	Program Counter:
(i)	SBIW Rdl,K	dl ∈ {24,26,28,30}, 0 ≤ K ≤ 63	PC ← PC + 1

16 bit Opcode:

1001	0111	KKdd	KKKK

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	⇔	⇔	⇔	⇔

S: N ⊕ V, For signed tests.

V: Rdh7 •$\overline{R15}$
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R15
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R15}$• $\overline{R14}$ •$\overline{R13}$ •$\overline{R12}$ •$\overline{R11}$• $\overline{R10}$• $\overline{R9}$• $\overline{R8}$• $\overline{R7}$• $\overline{R6}$ •$\overline{R5}$• $\overline{R4}$• $\overline{R3}$ •$\overline{R2}$• $\overline{R1}$• $\overline{R0}$
Set if the result is $0000; cleared otherwise.

C: R15• $\overline{Rdh7}$
Set if the absolute value of K is larger than the absolute value of Rd; cleared otherwise.

R (Result) equals Rdh:Rdl after the operation (Rdh7-Rdh0 = R15-R8, Rdl7-Rdl0=R7-R0).

Example:

```
sbiw    r24,1          ; Subtract 1 from r25:r24
sbiw    r28,63         ; Subtract 63 from the Y pointer(r29:r28)
```

Words: 1 (2 bytes)
Cycles: 2

SBR - Set Bits in Register

Description:

Sets specified bits in register Rd. Performs the logical ORI between the contents of register Rd and a constant mask K and places the result in the destination register Rd.

Operation:

(i) Rd ← Rd v K

	Syntax:	Operands:	Program Counter:
(i)	SBR Rd,K	16 ≤ d ≤ 31, 0 ≤ K ≤ 255	PC ← PC + 1

16 bit Opcode:

0110	KKKK	dddd	KKKK

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	**0**	⇔	⇔	-

S: N ⊕ V, For signed tests.

V: 0
Cleared

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}$• $\overline{R6}$ •$\overline{R5}$• $\overline{R4}$• $\overline{R3}$ •$\overline{R2}$• $\overline{R1}$• $\overline{R0}$
Set if the result is $00; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
sbr   r16,3          ; Set bits 0 and 1 in r16
sbr   r17,$F0        ; Set 4 MSB in r17
```

Words: 1 (2 bytes)
Cycles: 1

SBRC - Skip if Bit in Register is Cleared

Description:

This instruction tests a single bit in a register and skips the next instruction if the bit is cleared.

Operation:

(i) If Rr(b) = 0 then PC ← PC + 2 (or 3) else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	SBRC Rr,b	0 ≤ r ≤ 31, 0 ≤ b ≤ 7	PC ← PC + 1, If condition is false, no skip. PC ← PC + 2, If next instruction is one word. PC ← PC + 3, If next instruction is JMP or CALL

16 bit Opcode:

1111	110r	rrrr	Xbbb

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
sub   r0,r1        ; Subtract r1 from r0
sbrc  r0,7         ; Skip if bit 7 in r0 cleared
sub   r0,r1        ; Only executed if bit 7 in r0 not cleared
nop                ; Continue (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false (no skip)
2 if condition is true (skip is executed)

SBRS - Skip if Bit in Register is Set

Description:

This instruction tests a single bit in a register and skips the next instruction if the bit is set.

Operation:

(i) If Rr(b) = 1 then PC ← PC + 2 (or 3) else PC ← PC + 1

	Syntax:	Operands:	Program Counter:
(i)	SBRS Rr,b	0 ≤ r ≤ 31, 0 ≤ b ≤ 7	PC ← PC + 1, Condition false - no skip PC ← PC + 2, Skip a one word instruction PC ← PC + 3, Skip a JMP or a CALL

16 bit Opcode:

1111	111r	rrrr	Xbbb

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
sub     r0,r1       ; Subtract r1 from r0
sbrs    r0,7        ; Skip if bit 7 in r0 set
neg     r0          ; Only executed if bit 7 in r0 not set
nop                 ; Continue (do nothing)
```

Words: 1 (2 bytes)

Cycles: 1 if condition is false (no skip)
2 if condition is true (skip is executed)

SEC - Set Carry Flag

Description:
Sets the Carry flag (C) in SREG (status register).

Operation:
(i) C ← 1

	Syntax:	Operands:	Program Counter:
(i)	SEC	None	PC ← PC + 1

16 bit Opcode:

1001	0100	0000	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	**1**

C: 1
Carry flag set

Example:

```
sec               ; Set carry flag
adc       r0,r1   ; r0=r0+r1+1
```

Words: 1 (2 bytes)
Cycles: 1

SEH - Set Half Carry Flag

Description:
Sets the Half Carry (H) in SREG (status register).

Operation:
(i) H ← 1

	Syntax:	Operands:	Program Counter:
(i)	SEH	None	PC ← PC + 1

16 bit Opcode:

1001	0100	0101	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	1	-	-	-	-	-

H: 1
Half Carry flag set

Example:
```
seh          ; Set Half Carry flag
```

Words: 1 (2 bytes)
Cycles: 1

SEI - Set Global Interrupt Flag

Description:
Sets the Global Interrupt flag (I) in SREG (status register).

Operation:
(i) I ← 1

	Syntax:	Operands:	Program Counter:
(i)	SEI	None	PC ← PC + 1

16 bit Opcode:

1001	0100	0111	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
1	-	-	-	-	-	-	-

I: 1
Global Interrupt flag set

Example:
```
cli                 ; Disable interrupts
in    r13,$16       ; Read Port B
sei                 ; Enable interrupts
```

Words: 1 (2 bytes)
Cycles: 1

SEN - Set Negative Flag

Description:
Sets the Negative flag (N) in SREG (status register).

Operation:
(i) N ← 1

	Syntax:	Operands:	Program Counter:
(i)	SEN	None	PC ← PC + 1

16 bit Opcode:

1001	0100	0010	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	**1**	-	-

N: 1
Negative flag set

Example:

```
add   r2,r19       ; Add r19 to r2
sen                ; Set negative flag
```

Words: 1 (2 bytes)
Cycles: 1

SER - Set all bits in Register

Description:
Loads $FF directly to register Rd.

Operation:
(i) Rd ← $FF

	Syntax:	Operands:	Program Counter:
(i)	SER Rd	16 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode:

1110	1111	dddd	1111

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
clr   r16             ; Clear r16
ser   r17             ; Set r17
out   $18,r16         ; Write zeros to Port B
nop                   ; Delay (do nothing)
out   $18,r17         ; Write ones to Port B
```

Words: 1 (2 bytes)
Cycles: 1

SES - Set Signed Flag

Description:
Sets the Signed flag (S) in SREG (status register).

Operation:
(i) S ← 1

	Syntax:	Operands:	Program Counter:
(i)	SES	None	PC ← PC + 1

16 bit Opcode:

1001	0100	0100	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	**1**	-	-	-	-

S: 1
Signed flag set

Example:

```
add   r2,r19          ; Add r19 to r2
ses                   ; Set negative flag
```

Words: 1 (2 bytes)
Cycles: 1

SET - Set T Flag

Description:
Sets the T flag in SREG (status register).

Operation:
(i) T ← 1

	Syntax:	Operands:	Program Counter:
(i)	SET	None	PC ← PC + 1

16 bit Opcode:

1001	0100	0110	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	**1**	-	-	-	-	-	-

T: 1
T flag set

Example:
```
set                 ; Set T flag
```

Words: 1 (2 bytes)
Cycles: 1

SEV - Set Overflow Flag

Description:
Sets the Overflow flag (V) in SREG (status register).

Operation:
(i) V ← 1

	Syntax:	Operands:	Program Counter:
(i)	SEV	None	PC ← PC + 1

16 bit Opcode:

1001	0100	0011	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	**1**	-	-	-

V: 1
Overflow flag set

Example:

```
add   r2,r19         ; Add r19 to r2
sev                  ; Set overflow flag
```

Words: 1 (2 bytes)
Cycles: 1

SEZ - Set Zero Flag

Description:
Sets the Zero flag (Z) in SREG (status register).

Operation:
(i) Z ← 1

	Syntax:	Operands:	Program Counter:
(i)	SEZ	None	PC ← PC + 1

16 bit Opcode:

1001	0100	0001	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	**1**	-

Z: 1
Zero flag set

Example:
```
add   r2,r19          ; Add r19 to r2
sez                   ; Set zero flag
```

Words: 1 (2 bytes)
Cycles: 1

SLEEP

Description:

This instruction sets the circuit in sleep mode defined by the MCU control register. When an interrupt wakes up the MCU from a sleep state, the instruction following the SLEEP instruction will be executed before the interrupt handler is executed.

Operation:

Syntax:	Operands:	Program Counter:
SLEEP	None	PC ← PC + 1

16 bit Opcode:

1001	0101	100X	1000

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
mov     r0,r11          ; Copy r11 to r0
sleep                   ; Put MCU in sleep mode
```

Words: 1 (2 bytes)
Cycles: 1

ST - Store Indirect From Register to SRAM using Index X

Description:

Stores one byte indirect from Register to SRAM. The SRAM location is pointed to by the X (16 bits) pointer register in the register file. Memory access is limited to the current SRAM Page of 64K bytes. To access another SRAM page the RAMPX register in the I/O area has to be changed.

The X pointer register can either be left unchanged after the operation, or it can be incremented or decremented. These features are especially suited for stack pointer usage of the X pointer register.

Using the X pointer:

	Operation:		**Comment:**
(i)	(X) ← Rr		X: Unchanged
(ii)	(X) ← Rr	X ← X+1	X: Post incremented
(iii)	X ← X - 1	(X) ← Rr	X: Pre decremented

	Syntax:	**Operands:**	**Program Counter:**
(i)	ST X, Rr	0 ≤ r ≤ 31	PC ← PC + 1
(ii)	ST X+, Rr	0 ≤ r ≤ 31	PC ← PC + 1
(iii)	ST -X, Rr	0 ≤ r ≤ 31	PC ← PC + 1

16 bit Opcode :

(i)	1001	001r	rrrr	1100
(ii)	1001	001r	rrrr	1101
(iii)	1001	001r	rrrr	1110

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
clr     r27             ; Clear X high byte
ldi     r26,$20         ; Set X low byte to $20
st      X+,r0           ; Store r0 in SRAM loc. $20(X post inc)
st      X,r1            ; Store r1 in SRAM loc. $21
ldi     r26,$23         ; Set X low byte to $23
st      r2,X            ; Store r2 in SRAM loc. $23
st      r3,-X           ; Store r3 in SRAM loc. $22(X pre dec)
```

Words: 1 (2 bytes)
Cycles: 2

ST (STD) - Store Indirect From Register to SRAM using Index Y

Description:

Stores one byte indirect with or without displacement from Register to SRAM. The SRAM location is pointed to by the Y (16 bits) pointer register in the register file. Memory access is limited to the current SRAM Page of 64K bytes. To access another SRAM page the RAMPY register in the I/O area has to be changed.

The Y pointer register can either be left unchanged after the operation, or it can be incremented or decremented. These features are especially suited for stack pointer usage of the Y pointer register.

Using the Y pointer:

	Operation:		**Comment:**
(i)	(Y) ← Rr		Y: Unchanged
(ii)	(Y) ← Rr	Y ← Y+1	Y: Post incremented
(iii)	Y ← Y - 1	(Y) ← Rr	Y: Pre decremented
(iiii)	(Y+q) ← Rr		Y: Unchanged, q: Displacement

	Syntax:	**Operands:**	**Program Counter:**
(i)	ST Y, Rr	0 ≤ r ≤ 31	PC ← PC + 1
(ii)	ST Y+, Rr	0 ≤ r ≤ 31	PC ← PC + 1
(iii)	ST -Y, Rr	0 ≤ r ≤ 31	PC ← PC + 1
(iiii)	STD Y+q, Rr	0 ≤ r ≤ 31, 0 ≤ q ≤ 63	PC ← PC + 1

16 bit Opcode :

(i)	1000	001r	rrrr	1000
(ii)	1001	001r	rrrr	1001
(iii)	1001	001r	rrrr	1010
(iiii)	10q0	qq1r	rrrr	1qqq

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
clr     r29          ; Clear Y high byte
ldi     r28,$20      ; Set Y low byte to $20
st      Y+,r0        ; Store r0 in SRAM loc. $20(Y post inc)
st      Y,r1         ; Store r1 in SRAM loc. $21
ldi     r28,$23      ; Set Y low byte to $23
st      Y,r2         ; Store r2 in SRAM loc. $23
st      -Y,r3        ; Store r3 in SRAM loc. $22(Y pre dec)
std     Y+2,r4       ; Store r4 in SRAM loc. $24
```

Words: 1 (2 bytes)
Cycles: 2

ST (STD) - Store Indirect From Register to SRAM using Index Z

Description:

Stores one byte indirect with or without displacement from Register to SRAM. The SRAM location is pointed to by the Z (16 bits) pointer register in the register file. Memory access is limited to the current SRAM Page of 64K bytes. To access another SRAM page the RAMPZ register in the I/O area has to be changed.

The Z pointer register can either be left unchanged after the operation, or it can be incremented or decremented. These features are very suited for stack pointer usage of the Z pointer register, but because the Z pointer register can be used for indirect subroutine calls, indirect jumps and table lookup it is often more convenient to use the X or Y pointer as a dedicated stack pointer.

Using the Z pointer:

	Operation:		**Comment:**
(i)	(Z) ←Rr		Z: Unchanged
(ii)	(Z) ← Rr	Z ← Z+1	Z: Post incremented
(iii)	Z ← Z - 1	(Z) ← Rr	Z: Pre decremented
(iiii)	(Z+q) ← Rr		Z: Unchanged, q: Displacement

	Syntax:	**Operands:**	**Program Counter:**
(i)	ST Z, Rr	0 ≤ r ≤ 31	PC ← PC + 1
(ii)	ST Z+, Rr	0 ≤ r ≤ 31	PC ← PC + 1
(iii)	ST -Z, Rr	0 ≤ r ≤ 31	PC ← PC + 1
(iiii)	STD Z+q, Rr	0 ≤ r ≤ 31, 0 ≤ q ≤ 63	PC ← PC + 1

16 bit Opcode :

(i)	1000	001r	rrrr	0000
(ii)	1001	001r	rrrr	0001
(iii)	1001	001r	rrrr	0010
(iiii)	10q0	qq1r	rrrr	0qqq

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
clr     r31          ; Clear Z high byte
ldi     r30,$20      ; Set Z low byte to $20
st      Z+,r0        ; Store r0 in SRAM loc. $20(Z post inc)
st      Z,r1         ; Store r1 in SRAM loc. $21
ldi     r30,$23      ; Set Z low byte to $23
st      Z,r2         ; Store r2 in SRAM loc. $23
st      -Z,r3        ; Store r3 in SRAM loc. $22(Z pre dec)
std     Z+2,r4       ; Store r4 in SRAM loc. $24
```

Words: 1 (2 bytes)
Cycles: 2

STS - Store Direct to SRAM

Description:

Stores one byte from a Register to the SRAM. A 16-bit address must be supplied. Memory access is limited to the current SRAM Page of 64K bytes. The SDS instruction uses the RAMPZ register to access memory above 64K bytes.

Operation:

(i) (k) ← Rr

	Syntax:	Operands:	Program Counter:
(i)	STS k,Rr	0 ≤ r ≤ 31, 0 ≤ k ≤ 65535	PC ← PC + 2

32 bit Opcode:

1001	001d	dddd	0000
kkkk	kkkk	kkkk	kkkk

Status Register (SREG) and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
lds     r2,$FF00        ; Load r2 with the contents of SRAM location $FF00
add     r2,r1           ; add r1 to r2
sts     $FF00,r2        ; Write back
```

Words: 2 (4 bytes)
Cycles: 3

SUB - Subtract without Carry

Description:
Subtracts two registers and places the result in the destination register Rd.

Operation:
(i) Rd ← Rd - Rr

	Syntax:	Operands:	Program Counter:
(i)	SUB Rd,Rr	$0 \le d \le 31, 0 \le r \le 31$	PC ← PC + 1

16 bit Opcode:

0001	10rd	dddd	rrrr

Status Register and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: $\overline{Rd3}\bullet Rr3 + Rr3 \bullet R3 + R3 \bullet \overline{Rd3}$
Set if there was a borrow from bit 3; cleared otherwise

S: $N \oplus V$, For signed tests.

V: $Rd7 \bullet \overline{Rr7} \bullet \overline{R7} + \overline{Rd7} \bullet Rr7 \bullet R7$
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; cleared otherwise.

C: $\overline{Rd7} \bullet Rr7 + Rr7 \bullet R7 + R7 \bullet \overline{Rd7}$
Set if the absolute value of the contents of Rr is larger than the absolute value of Rd; cleared otherwise.

R (Result) equals Rd after the operation.

Example:
```
            sub     r13,r12      ; Subtract r12 from r13
            brne    noteq        ; Branch if r12<>r13
            ...
noteq:      nop                  ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

SUBI - Subtract Immediate

Description:

Subtracts a register and a constant and places the result in the destination register Rd. This instruction is working on Register R16 to R31 and is very well suited for operations on the X, Y and Z pointers.

Operation:

(i) Rd ← Rd - K

	Syntax:	Operands:	Program Counter:
(i)	SUBI Rd,K	16 ≤ d ≤ 31, 0 ≤ K ≤ 255	PC ← PC + 1

16 bit Opcode:

0101	KKKK	dddd	KKKK

Status Register and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	⇔	⇔	⇔	⇔	⇔	⇔

H: $\overline{Rd3}$• K3+K3 •R3 +R3 •$\overline{Rd3}$
Set if there was a borrow from bit 3; cleared otherwise

S: N ⊕ V, For signed tests.

V: Rd7• $\overline{K7}$ •$\overline{R7}$ +$\overline{Rd7}$• K7 •R7
Set if two's complement overflow resulted from the operation; cleared otherwise.

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7}$• $\overline{R6}$ •$\overline{R5}$• $\overline{R4}$• $\overline{R3}$ •$\overline{R2}$• $\overline{R1}$• $\overline{R0}$
Set if the result is $00; cleared otherwise.

C: $\overline{Rd7}$• K7 +K7 •R7 +R7• $\overline{Rd7}$
Set if the absolute value of K is larger than the absolute value of Rd; cleared otherwise.

R (Result) equals Rd after the operation.

Example:

```
            subir22,$11      ; Subtract $11 from r22
            brnenoteq        ; Branch if r22<>$11
            ...
   noteq:   nop              ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

SWAP - Swap Nibbles

Description:
Swaps high and low nibbles in a register.

Operation:
(i) R(7-4) ← Rd(3-0), R(3-0) ← Rd(7-4)

	Syntax:	Operands:	Program Counter:
(i)	SWAP Rd	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode:

1001	010d	dddd	0010

Status Register and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

R (Result) equals Rd after the operation.

Example:

```
inc     r1              ; Increment r1
swap    r1              ; Swap high and low nibble of r1
inc     r1              ; Increment high nibble of r1
swap    r1              ; Swap back
```

Words: 1 (2 bytes)
Cycles: 1

TST - Test for Zero or Minus

Description:

Tests if a register is zero or negative. Performs a logical AND between a register and itself. The register will remain unchanged.

Operation:

(i) Rd ← Rd • Rd

	Syntax:	Operands:	Program Counter:
(i)	TST Rd	0 ≤ d ≤ 31	PC ← PC + 1

16 bit Opcode:

0010	00dd	dddd	dddd

Status Register and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	⇔	**0**	⇔	⇔	-

S: N ⊕ V, For signed tests.

V: 0
Cleared

N: R7
Set if MSB of the result is set; cleared otherwise.

Z: $\overline{R7} \bullet \overline{R6} \bullet \overline{R5} \bullet \overline{R4} \bullet \overline{R3} \bullet \overline{R2} \bullet \overline{R1} \bullet \overline{R0}$
Set if the result is $00; cleared otherwise.

R (Result) equals Rd.

Example:

```
            tst   r0         ; Test r0
            breq  zero       ; Branch if r0=0
            ...
zero:       nop              ; Branch destination (do nothing)
```

Words: 1 (2 bytes)
Cycles: 1

WDR - Watchdog Reset

Description:

This instruction resets the Watchdog Timer. This instruction must be executed within a limited time given by the WD prescaler. See the Watchdog Timer hardware specification.

Operation:

(i) WD timer restart.

	Syntax:	Operands:	Program Counter:
(i)	WDR	None	PC ← PC + 1

16 bit Opcode:

1001	0101	101X	1000

Status Register and Boolean Formulae:

I	T	H	S	V	N	Z	C
-	-	-	-	-	-	-	-

Example:

```
wdr             ; Reset watchdog timer
```

Words: 1 (2 bytes)
Cycles: 1